

Lecture

- In the lecture on April 24 we discussed the rest of Part Four of the book (mainly Chapter).

Exercises

Note, as usual, that you find even more exercises including solutions here :

<http://codex.cs.yale.edu/avi/os-book/OS9/practice-exer-dir/index.html>

Prepare for the Tutorial Session in week 18, 2018: all exercises that have not been discussed so far as well as the following:

- 12.1 None of the disk-scheduling disciplines, except FCFS, is truly fair (starvation may occur).
- a. Explain why this assertion is true.
 - b. Describe a way to modify algorithms such as SCAN to ensure fairness.
 - c. Explain why fairness is an important goal in a time-sharing system.
 - d. Give three or more examples of circumstances in which it is important that the operating system be unfair in serving I/O requests.
- 12.2 Explain why SSDs often use a FCFS disk scheduling algorithm.
- 12.3 Suppose that a disk drive has 5000 cylinders, numbered 0 to 4999. The drive is currently serving a request at cylinder 2150, and the previous request was at cylinder 1805. The queue of pending requests, in FIFO order, is:

2069, 1212, 2296, 2800, 544, 1618, 356, 1523, 4965, 3681.

Starting from the current head position, what is the total distance (in cylinders) that the disk arm moves to satisfy all the pending requests for each of the following disk-scheduling algorithms?

- a. FCFS
- b. SSTF
- c. SCAN
- d. LOOK
- e. C-SCAN
- f. C-LOOK

- 12.4 Elementary physics states that when an object is subjected to a constant acceleration a , the relationship between distance d and time t is given by $d = 1/2at^2$. Suppose that, during a seek, the disk in Exercise 12.3 accelerates the disk arm at a constant rate for the first half of the seek, then decelerates the disk arm at the same rate for the second half of the seek. Assume that the disk can perform a seek to an adjacent cylinder in 1 millisecond and a full-stroke seek over all 5000 cylinders in 18 milliseconds.
- The distance of a seek is the number of cylinders that the head moves. Explain why the seek time is proportional to the square root of the seek distance.
 - Write an equation for the seek time as a function of the seek distance. This equation should be of the form $t = x + y\sqrt{L}$, where t is the time in milliseconds and L is the seek distance in cylinders.
 - Calculate the total seek time for each of the schedules in Exercise 12.3. Determine which schedule is the fastest (has the smallest total seek time).
 - The percentage speedup is the time saved divided by the original time. What is the percentage speedup of the fastest schedule over FCFS?
- 12.6 Describe some advantages and disadvantages of using SSDs as a caching tier and as a disk drive replacement compared to a system with just magnetic disks.
- 12.7 Compare the performance of C-SCAN and SCAN scheduling, assuming a uniform distribution of requests. Consider the average response time (the time between the arrival of a request and the completion of that request's service), the variation in response time, and the effective bandwidth. How does performance depend on the relative sizes of seek time and rotational latency?
- 12.8 Requests are not usually uniformly distributed. For example, we can expect a cylinder containing the file-system metadata to be accessed more frequently than a cylinder containing only files. Suppose you know that 50 percent of the requests are for a small, fixed number of cylinders.
- Would any of the scheduling algorithms discussed in this chapter be particularly good for this case? Explain your answer.
 - Propose a disk-scheduling algorithm that gives even better performance by taking advantage of this "hot spot" on the disk.
- 12.9 Consider a RAID Level 5 organization comprising five disks, with the parity for sets of four blocks on four disks stored on the fifth disk. How many blocks are accessed in order to perform the following?
- A write of one block of data
 - A write of seven continuous blocks of data

- 12.10 Compare the throughput achieved by a RAID Level 5 organization with that achieved by a RAID Level 1 organization for the following:
- Read operations on single blocks
 - Read operations on multiple contiguous blocks
- 12.11 Compare the performance of write operations achieved by a RAID Level 5 organization with that achieved by a RAID Level 1 organization.
- 12.12 Assume that you have a mixed configuration comprising disks organized as RAID Level 1 and as RAID Level 5 disks. Assume that the system has flexibility in deciding which disk organization to use for storing a particular file. Which files should be stored in the RAID Level 1 disks and which in the RAID Level 5 disks in order to optimize performance?
- 12.13 The reliability of a hard-disk drive is typically described in terms of a quantity called mean time between failures (MTBF). Although this quantity is called a “time”, the MTBF actually is measured in drivehours per failure.
- If a system contains 1000 drives, each of which has a 750,000-hour MTBF, which of the following best describes how often a drive failure will occur in that disk farm: once per thousand years, once per century, once per decade, once per year, once per month, once per week, once per day, once per hour, once per minute, or once per second?
 - Mortality statistics indicate that, on the average, a U.S. resident has about 1 chance in 1000 of dying between ages 20 and 21 years. Deduce the MTBF hours for 20 year olds. Convert this figure from hours to years. What does this MTBF tell you about the expected lifetime of a 20 year old?
 - The manufacturer guarantees a 1-million-hour MTBF for a certain model of disk drive. What can you conclude about the number of years for which one of these drives is under warranty?
- 12.14 Discuss the relative advantages and disadvantages of sector sparing and sector slipping.
- 12.15 Discuss the reasons why the operating system might require accurate information on how blocks are stored on a disk. How could the operating system improve file system performance with this knowledge?
- 13.1 When multiple interrupts from different devices appear at about the same time, a priority scheme could be used to determine the order in which the interrupts would be serviced. Discuss what issues need to be considered in assigning priorities to different interrupts.
- 13.2 What are the advantages and disadvantages of supporting memory-mapped I/O to device-control registers?

13.3 Consider the following I/O scenarios on a single-user PC.

- a. A mouse used with a graphical user interface
- b. A tape drive on a multitasking operating system (assume no device preallocation is available)
- c. A disk drive containing user files
- d. A graphics card with direct bus connection, accessible through memory-mapped I/O

For each of these I/O scenarios, would you design the operating system to use buffering, spooling, caching, or a combination? Would you use polled I/O, or interrupt-driven I/O? Give reasons for your choices.

13.4 In most multiprogrammed systems, user programs access memory through virtual addresses, while the operating system uses raw physical addresses to access memory. What are the implications of this design on the initiation of I/O operations by the user program and their execution by the operating system?

13.5 What are the various kinds of performance overheads associated with servicing an interrupt?

13.6 Describe three circumstances under which blocking I/O should be used. Describe three circumstances under which non-blocking I/O should be used. Why not just implement non-blocking I/O and have processes busy-wait until their device is ready?

13.7 Typically, at the completion of a device I/O, a single interrupt is raised and appropriately handled by the host processor. In certain settings, however, the code that is to be executed at the completion of the I/O can be broken into two separate pieces. The first piece executes immediately after the I/O completes and schedules a second interrupt for the remaining piece of code to be executed at a later time. What is the purpose of using this strategy in the design of interrupt handlers?

13.8 Some DMA controllers support direct virtual memory access, where the targets of I/O operations are specified as virtual addresses and a translation from virtual to physical address is performed during the DMA. How does this design complicate the design of the DMA controller? What are the advantages of providing such a functionality?