Motivation

- A marked net $\langle N, M_0 \rangle$ with N = (P, T, Pre, Post) specifies:
 - an initial marking (i.e., state) *M*₀;
 - the rules of evolution.

No explicit enumeration of:

net language, i.e., the set of sequences of transitions that can fire:

 $L(N, M_0) = \{ \sigma \in T^* \mid M_0[\sigma \rangle \};$

reachability set, i.e., the set of reachable markings:

 $R(N, M_0) = \{ M \in \mathbb{N}^{|P|} \mid (\exists \sigma \in L(N, M_0)) \ M_0[\sigma \rangle M \}.$

The information on reachable markings and firing sequences is useful to determine if the net has given properties.

Reachability graph

The reachability graph of a marked net $\langle N, M_0 \rangle$ is an automaton $\mathcal{G} = (X, E, \delta, x_0)$ where:

- X = R(N, M₀), i.e., the states of the automaton are the reachable markings;
- E = T, i.e., the events in the alphabet are the transitions of the net;
- for any two reachable markings M, M':

 $\delta(M,t) = M' \quad \Longleftrightarrow \quad M[t\rangle M',$

i.e., there exists arc labeled t from M to M' on the automaton iff marking M' is reachable from M firing transition t;

• $x_0 = M_0$, i.e., the initial state of the automaton is the initial marking.

It can be constructed only if the reachability set if finite, i.e., if the net is bounded.