## Mathematical Optimization



Marco Chiarandini


Institut for Matematik og Datalogi (IMADA)

October 25, 2018
MatØk Studiepratik

## Operations Research

Operation Research (aka, Management Science, Analytics):

- the discipline that uses a scientific approach to decision making.
- It seeks to determine how best to design and operate a system, usually under conditions requiring the allocation of scarce resources, by means of mathematics and computer science.
- Quantitative methods for planning and analysis.

Applications:

- Transport
- Supply Chains
- Sport
- Finance
- Government
- Manufacturing


## Today’s Objectives

- Convert a problem from ordinary language into mathematical language
- Solve geometrically a system of linear inequalities and connect the solutions to the real world problem
- Distinguish Linear vs Non-linear, Continuous vs Integer problems
- Solve numerically the problems in Google Sheets and Microsoft Excel
- Applications: production planning, diet planning, budget allocation
$\rightsquigarrow$ Transmit to you my fascination for Mathematical Optimization

1. Production Planning
2. Diet Problem
3. Budget Allocation
4. Summary

## Outline

1. Production Planning
2. Diet Problem
3. Budget Allocation
4. Summary

## Production Planning

Suppose a company produces only tables and chairs.
A table is made of 2 large Lego pieces and 2 small pieces, while a chair is made of 1 large and 2 small pieces.
The resources available are 8 small and 6 large pieces.


A table


A chair



The profit for a table is 1600 dkk and for a chair 1000 dkk . What product mix maximizes the company's profile using the available resources?

|  | Tables | Chairs | Capacity |
| ---: | :---: | :---: | :---: |
| Small Pieces | 2 | 2 | 8 |
| Large Pieces | 2 | 1 | 6 |
| Profit | 16 | 10 |  |

Decision Variables

$$
\begin{aligned}
& x_{1} \geq 0 \text { units of small pieces } \\
& x_{2} \geq 0 \text { units of large pieces }
\end{aligned}
$$

## Object Function

$$
\max 16 x_{1}+10 x_{2} \text { maximize profit }
$$

Constraints

$$
\begin{aligned}
& 2 x_{1}+2 x_{2} \leq 8 \text { small pieces capacity } \\
& 2 x_{1}+x_{2} \leq 6 \text { large pieces capacity }
\end{aligned}
$$

Materials $A$ and $B$
Products 1 and 2

$$
\begin{aligned}
\max 16 x_{1}+10 x_{2} & \\
2 x_{1}+2 x_{2} & \leq 8 \\
2 x_{1}+x_{2} & \leq 6 \\
x_{1} & \geq 0 \\
x_{2} & \geq 0
\end{aligned}
$$

Graphical Representation:


## Resource Allocation - General Model

Managing a production facility


$$
\begin{aligned}
& \max c_{1} x_{1}+c_{2} x_{2}+c_{3} x_{3}+\ldots+c_{n} x_{n}=z \\
& \text { s.t. } a_{11} x_{1}+a_{12} x_{2}+a_{13} x_{3}+\ldots+a_{1 n} x_{n} \leq b_{1} \\
& a_{21} x_{1}+a_{22} x_{2}+a_{23} x_{3}+\ldots+a_{2 n} x_{n} \leq b_{2} \\
& a_{m 1} x_{1}+a_{m 2} x_{2}+a_{m 3} x_{3}+\ldots+a_{m n} x_{n} \leq b_{m} \\
& x_{1}, x_{2}, \ldots, x_{n} \geq 0
\end{aligned}
$$

$$
\begin{aligned}
\max \quad \sum_{j=1}^{n} c_{j} x_{j} & \\
\sum_{j=1}^{n} a_{i j} x_{j} & \leq b_{i}, \quad i=1, \ldots, m \\
x_{j} & \geq 0, \quad j=1, \ldots, n
\end{aligned}
$$

$$
\begin{aligned}
c^{T}=\left[\begin{array}{llll}
c_{1} & c_{2} & \ldots & c_{n}
\end{array}\right] \\
A=\left[\begin{array}{cccc}
a_{11} & a_{12} & \ldots & a_{1 n} \\
a_{21} & a_{22} & \ldots & a_{2 n} \\
\vdots & & & \\
a_{31} & a_{32} & \ldots & a_{m n}
\end{array}\right], x=\left[\begin{array}{c}
x_{1} \\
x_{2} \\
\vdots \\
\vdots \\
x_{n}
\end{array}\right], b=c^{T} x \\
x \geq 0
\end{aligned}
$$

$$
\begin{aligned}
& \max \quad c_{1} x_{1}+c_{2} x_{2}+c_{3} x_{3}+\ldots+c_{n} x_{n}=z \\
& \text { s.t. } a_{11} x_{1}+a_{12} x_{2}+a_{13} x_{3}+\ldots+a_{1 n} x_{n} \leq b_{1} \\
& a_{21} x_{1}+a_{22} x_{2}+a_{23} x_{3}+\ldots+a_{2 n} x_{n} \leq b_{2} \\
& \begin{aligned}
a_{m 1} x_{1}+a_{m 2} x_{2}+a_{m 3} x_{3}+\ldots+a_{m n} x_{n} & \leq b_{m} \\
x_{1}, x_{2}, \ldots, x_{n} & \geq 0
\end{aligned}
\end{aligned}
$$

## Vector and Matrices in Excel

$$
\sum_{j=1}^{n} c_{j}=c_{1}+c_{2}+\ldots+c_{n}
$$

SUM(B5: B14)

## Scalar product

$$
\begin{aligned}
\vec{u} \cdot \vec{v} & =u_{1} v_{1}+u_{2} v_{2}+\ldots+u_{n} v_{n} \\
& =\sum_{j=1}^{n} u_{j} v_{j}
\end{aligned}
$$

SUMPRODUCT(B5: B14, C5: C : 14)

$$
\begin{aligned}
& \max \quad \sum_{j=1}^{n} c_{j} x_{j} \\
& \sum_{j=1}^{n} a_{i j} x_{j} \leq b_{i}, \quad i=1, \ldots, m \\
& x_{j} \geq 0, \quad j=1, \ldots, n
\end{aligned}
$$

$\max c^{\top} x$

$$
A x \leq b
$$

$$
x \geq 0
$$

$$
x \in \mathbb{R}^{n}, c \in \mathbb{R}^{n}, A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^{m}
$$

$$
\begin{aligned}
\max 16 x_{1}+10 x_{2} & \\
2 x_{1}+2 x_{2} & \leq 8 \\
2 x_{1}+x_{2} & \leq 6 \\
x_{1}, x_{2} & \geq 0
\end{aligned}
$$

$$
\max \left[\begin{array}{ll}
16 & 10
\end{array}\right]\left[\begin{array}{l}
x_{1} \\
x_{2}
\end{array}\right]
$$

$$
\left[\begin{array}{ll}
2 & 2 \\
2 & 1
\end{array}\right]\left[\begin{array}{l}
x_{1} \\
x_{2}
\end{array}\right] \leq\left[\begin{array}{l}
8 \\
6
\end{array}\right]
$$

$$
x_{1}, x_{2} \geq 0
$$

## 1．Production Planning

2．Diet Problem

## 3．Budget Allocation

## 4．Summary

## The Diet Problem (Blending Problems)

- Select a set of foods that will satisfy a set of daily nutritional requirement at minimum cost.
- Motivated in the 1930s and 1940s by US army.
- Formulated as a linear programming problem by George Stigler
- First linear programming problem
- (programming intended as planning not computer code)

min cost/weight
subject to nutrition requirements:
eat enough but not too much of Vitamin A eat enough but not too much of Sodium eat enough but not too much of Calories


## The Diet Problem

Suppose there are:

- 3 foods available, corn, milk, and bread,
- there are restrictions on the number of calories (between 2000 and 2250) and the amount of Vitamin A (between 5000 and 50,000 )

| Food | Corn | $2 \%$ Milk | Wheat bread |
| :--- | ---: | ---: | ---: |
| Vitamin A | 107 | 500 | 0 |
| Calories | 72 | 121 | 65 |
| Cost per serving | $\$ 0.18$ | $\$ 0.23$ | $\$ 0.05$ |

```
Parameters (given data)
    F = set of foods
    N = set of nutrients
    aij = amount of nutrient j in food i,\foralli\inF,\forallj\inN
    ci}=\mathrm{ cost per serving of food i,}\foralli\in
    F
    F}\mp@subsup{F}{maxi}{}=\mathrm{ maximum allowable number of servings of food i,}\foralli\in
    N Ninj }=\mathrm{ minimum required level of nutrient j,},\forallj\in
    Nmaxj = maximum allowable level of nutrient j,\forallj\inN
```

Decision Variables
$x_{i}=$ number of servings of food $i$ to purchase/consume, $\forall i \in F$

## The Mathematical Model

Objective Function: Minimize the total cost of the food

$$
\operatorname{Minimize} \sum_{i \in F} c_{i} x_{i}
$$

Constraint Set 1: For each nutrient $j \in N$, at least meet the minimum required level

$$
\sum_{i \in F} a_{i j} x_{i} \geq N_{\operatorname{minj}}, \quad \forall j \in N
$$

Constraint Set 2: For each nutrient $j \in N$, do not exceed the maximum allowable level.

$$
\sum_{i \in F} a_{i j} x_{i} \leq N_{\max j}, \quad \forall j \in N
$$

Constraint Set 3: For each food $i \in F$, select at least the minimum required number of servings

$$
x_{i} \geq F_{\operatorname{mini}}, \quad \forall i \in F
$$

Constraint Set 4: For each food $i \in F$, do not exceed the maximum allowable number of servings.

$$
x_{i} \leq F_{\operatorname{maxi}}, \quad \forall i \in F
$$

system of equalities and inequalities

$$
\begin{aligned}
\min \sum_{i \in F} c_{i} x_{i} & \\
\sum_{i \in F} a_{i j} x_{i} \geq N_{\operatorname{minj} j}, & \forall j \in N \\
\sum_{i \in F} a_{i j} x_{i} \leq N_{\max j}, & \forall j \in N \\
x_{i} \geq F_{\min i}, & \forall i \in F \\
x_{i} \leq F_{\max i}, & \forall i \in F
\end{aligned}
$$

## The History of Stigler's Diet Problem

- The linear program consisted of 9 equations in 77 variables
- Stigler, guessed an optimal solution using a heuristic method
- In 1947, the National Bureau of Standards used the newly developed simplex method to solve Stigler's model.
It took 9 clerks using hand-operated desk calculators 120 man days to solve for the optimal solution


## Geometrical Interpretation

Geometrically the feasibility region of a linear programming problem with 3 variables is a polyhedron.


The generalization of a polyhedron in $n$ dimensions is called polytope.

## Growth Functions



NP-hard problems: bad if we have to solve them, good for cryptology

# 1. Production Planning 

2. Diet Problem
3. Budget Allocation
4. Summary

## Budget Allocation

- A company has six different opportunities to invest money.
- Each opportunity requires a certain investment over a period of 6 years or less.

| Expected <br> Investment Cash <br> Flows and Net <br> Present Value |  |  |  |  |  |  |  |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
|  | Opp. 1 | Opp. 2 | Opp. 3 | Opp. 4 | Opp. 5 | Opp. 6 |  |
| Year 1 | $-\$ 5.00$ | $-\$ 9.00$ | $-\$ 12.00$ | $-\$ 7.00$ | $-\$ 20.00$ | $-\$ 18.00$ |  |
| Year 2 | $-\$ 6.00$ | $-\$ 6.00$ | $-\$ 10.00$ | $-\$ 5.00$ | $\$ 6.00$ | $-\$ 15.00$ |  |
| Year 3 | $-\$ 16.00$ | $\$ 6.10$ | $-\$ 5.00$ | $-\$ 20.00$ | $\$ 6.00$ | $-\$ 10.00$ |  |
| Year 4 | $\$ 12.00$ | $\$ 4.00$ | $-\$ 5.00$ | $-\$ 10.00$ | $\$ 6.00$ | $-\$ 10.00$ |  |
| Year 5 | $\$ 14.00$ | $\$ 5.00$ | $\$ 25.00$ | $-\$ 15.00$ | $\$ 6.00$ | $\$ 35.00$ |  |
| Year 6 | $\$ 15.00$ | $\$ 5.00$ | $\$ 15.00$ | $\$ 75.00$ | $\$ 6.00$ | $\$ 35.00$ |  |
| NPV | $\$ 8.01$ | $\$ 2.20$ | $\$ 1.85$ | $\$ 7.51$ | $\$ 5.69$ | $\$ 5.93$ |  |

- The company wants to invest in those opportunities that maximize the combined Net Present Value (NPV).
- It also has an investment budget that needs to be met for each year.


## Net Present Value

- $P$ : value of the original payment presently due
- the debtor wants to delay the payment for $t$ years,
- let $r$ be the market rate of return on a similar investment asset
- the future value of $P$ is $F=P(1+r)^{t}$

Viceversa, consider the task of finding:

- the present value $P$ of $\$ 100$ that will be received in five years, or equivalently,
- which amount of money today will grow to $\$ 100$ in five years when subject to a constant discount rate.

Assuming a 5\% per year interest rate, it follows that

$$
P=\frac{F}{(1+r)^{t}}=\frac{\$ 100}{(1+0.05)^{5}}=\$ 78.35 .
$$

## Budget Allocation

Net Present Value calculation:
for each opportunity we calculate the NPV at time zero (the time of decision) as:

$$
P_{0}=\sum_{t=1}^{5} \frac{F_{t}}{(1+0.05)^{5}}
$$

| Expected <br> Investment Cash <br> Flows and Net <br> Present Value |  |  |  |  |  |  |  |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: |

## Budget Allocation - Mathematical Model

- Let $B_{t}$ be the budget available for investments during the years $t=1 . .5$.
- Let $a_{t j}$ be the cash flow for opportunity $j$ and $c_{j}$ its NPV
- Task: choose a set of opportunities such that the budget is never exceeded and the expected return is maximized. Consider both the case of indivisible and divisible opportunities.

Variables $x_{j}=1$ if opportunity $j$ is selected and $x_{j}=0$ otherwise, $j=1 . .6$
Objective

$$
\max \sum_{j=1}^{6} c_{j} x_{j}
$$

Constraints

$$
\sum_{j=1}^{6} a_{t j} x_{j}+B_{t} \geq 0 \quad \forall t=1 . .5
$$

1．Production Planning
2．Diet Problem
3．Budget Allocation
4．Summary


## Mathematical Modeling

- Find out exactly what the decision makers need to know:
- which investment?
- which product mix?
- which job $j$ should a person $i$ do?
- Define Decision Variables of suitable type (continuous, integer valued, binary) corresponding to the needs
- Formulate Objective Function computing the benefit/cost
- Formulate mathematical Constraints indicating the interplay between the different variables.

Recognize linear and non linear functions and continuous and integer variables.

- Geometrical interpretation of the simplex method
- Touched computational issues
- Computer carries out the operations, hence programming needed
- Practical experience with Spreadsheets


## Summary

1. Production Planning
2. Diet Problem
3. Budget Allocation
4. Summary
