

DM826 – Spring 2012  
Modeling and Solving Constrained Optimization Problems

Exercises  
Set Variables  
SONET Problem

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*[Partly based on slides by Stefano Gualandi, Politecnico di Milano]*

# Sonet problem

Optical fiber network design

## Sonet problem

**Input:** weighted undirected demand graph  $G = (N, E; d)$ , where each node  $u \in N$  represents a **client** and weighted edges  $(u, v) \in E$  correspond to **traffic demands** of a pair of clients.

Two nodes can communicate, only if they join the same **ring**; nodes may join more than one ring. We must respect:

- maximum number of rings  $r$
- maximum number of clients per ring  $a$
- maximum bandwidth capacity of each ring  $c$

**Task:** find a topology that minimizes the sum, over all rings, of the number of nodes that join each ring while clients' traffic demands are met.

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A solution of the SONET problem is an assignment of rings to nodes and of capacity to demands such that

1. all demands of each client pairs are satisfied;
2. the ring traffic does not exceed the bandwidth capacity;
3. at most  $r$  rings are used;
4. at most  $a$  ADMs on each ring;
5. the total number of ADMs used is minimized.

## Sonet: variables

- Set variable  $X_i$  represents the set of nodes assigned to ring  $i$
- Set variable  $Y_u$  represents the set of rings assigned to node  $u$
- Integer variable  $Z_{ie}$  represents the amount of bandwidth assigned to demand pair  $e$  on ring  $i$ .

## Sonnet: model

$$\begin{aligned} \min \quad & \sum_{i \in R} |X_i| \\ \text{s.t.} \quad & |Y_u \cap Y_v| \geq 1, & \forall (u, v) \in E, \\ & Z_{ie} \in \{0, d(e)\}, & \forall e \in E, \\ & Z_{i,(u,v)} > 0 \iff i \in (Y_u \cap Y_v), & \forall i \in R, (u, v) \in E, \\ & u \in X_i \iff i \in Y_u, & \forall i \in R, u \in N, \\ & |X_i| \leq a, & \forall i \in R \\ & \sum_{e \in E} Z_{ie} \leq c, & \forall i \in R. \\ & X_i \preceq X_j, & \forall i, j \in R : i < j. \end{aligned}$$

```

from numpy import *
from gcode import *
Rings = range(4) # upper bound for amount of rings
Nodes = range(5) # amount of clients
demand = array([[0,1,0,1,1],
                [1,0,1,0,0],
                [0,1,0,0,1],
                [1,0,0,0,0],
                [1,0,1,0,0]])
capacity = [3,2,2,3] # capacity in nodes of possible rings

X = map(lambda r: m.setvar(intset(),0,len(Nodes),0,capacity[r]),Rings) #nodes for r
Y = m.setvars(len(Nodes),intset(),0,len(Rings),0,len(Rings)) # rings for u

# at least two nodes in each ring
for r in Rings:
    cardinality(X[r], IRT_NQ, 1) # implied constraint

for (n1,n2) in combinations(Nodes,2):
    IntVar z(intset(),0,4,1,len(Rings));
    if demand[n1,n2]==1:
        rel(Y[n1], SOT_INTER, Y[n2], SRT_SUP, z)

channel(X,Y)

IntVarArray z(len(Rings), )
for r in Rings:
    cardinality(X[r],z[r])

IntVar adm(0,len(Rings)*len(Nodes)
linear(z, IRT_EQ, adm)

```

Once a variable  $X[i]$  has been chosen, we first try to include the node that has the most communication with the nodes already placed in ring  $i$ ,

```
while (!and(i in Rings)(X[i].bound())) {
  selectMin (i in Rings: !X[i].bound())(X[i].getCardinalityVariable().getSize()) {
    set{int} S = X[i].getPossibleSet();
    set{int} R = X[i].getRequiredSet();
    Solver<CP> cp = X[i].getSolver();
    selectMax (e in S: !R.contains(e))(sum(n in R)(demand[e,n])) {
      try<cp> cp.requires(nodesInRing[i],e); | cp.excludes(nodesInRing[i],e);
    }
  }
}
```

# References