DM841 Discrete Optimization

#### Lecture 8

### Stochastic Local Search and Metaheuristics Simulated Annealing, Tabu Search, Variable Neighborhood Search

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- 1. Vertex Coloring
- 2. Stochastic Local Search
  - Randomized Iterative Improvement

#### 1. Trajectory Based Metaheuristics

Stochastic Local Search Simulated Annealing Iterated Local Search Tabu Search Variable Neighborhood Search Guided Local Search

#### 1. Trajectory Based Metaheuristics Stochastic Local Search

Simulated Annealing Iterated Local Search Tabu Search Variable Neighborhood Search Guided Local Search

# Min-Conflict Heuristic

For a general CSP:

```
procedure MCH (P, maxSteps)
```

**input:** *CSP instance P, positive integer maxSteps* **output:** *solution of P or* "no solution found"

*a* := randomly chosen assignment of the variables in *P*;

for step := 1 to maxSteps do

if a satisfies all constraints of P then return a end

- x := randomly selected variable from conflict set K(a);
- v := randomly selected value from the domain of x such that
  - setting x to v minimises the number of unsatisfied constraints;

a := a with x set to v;

end

return "no solution found" end MCH

# Min-Conflict Heuristic

#### On the N-queens problem:

```
import cotls:
int n = 16:
range Size = 1..n;
UniformDistribution distr(Size);
Solver<LS> m();
var{int} queen[Size](m,Size) := distr.get();
ConstraintSystem<LS> S(m);
S.post(alldifferent(queen));
S.post(alldifferent(all(i in Size) gueen[i] + i));
S.post(alldifferent(all(i in Size) queen[i] -i);
m.close();
int it = 0:
while (S.violations() > 0 \&\& it < 50 * n) {
  select(q in Size : S.violations(queen[q])>0) {
    selectMin(v in Size)(S.getAssignDelta(queen[q],v)) {
      queen[q] := v;
      cout<<"chng @ "<<it<<": queen["<<q<<"] := "<<v<<" viol: "<<S.violations()
            <<endl:
    it = it + 1:
cout << queen << endl:
```

### Min-Conflict + Random Walk for SAT

Trajectory Based Metaheuris Population Based Metaheuris

procedure WalkSAT (F, maxTries, maxSteps, slc)

input: CNF formula F, positive integers maxTries and maxSteps, heuristic function slc

output: model of F or 'no solution found'

for try := 1 to maxTries do

a := randomly chosen assignment of the variables in formula F;

for step := 1 to maxSteps do

if a satisfies F then return a end

c := randomly selected clause unsatisfied under a;

x := variable selected from *c* according to heuristic function *slc*;

a := a with x flipped;

#### end

end

return 'no solution found' end WalkSAT

Example of *slc* heuristic: with prob. *wp* select a random move, with prob. 1 - wp select the best

# Probabilistic Iterative Improv.

**Key idea:** Accept worsening steps with probability that depends on respective deterioration in evaluation function value: bigger deterioration  $\cong$  smaller probability

Realization:

► Function p(f, s): determines probability distribution over neighbors of s based on their values under evaluation function f.

```
• Let step(s, s') := p(f, s, s').
```

Note:

- Behavior of PII crucially depends on choice of p.
- ▶ II and RII are special cases of PII.

#### Example: Metropolis PII for the TSP

- **Search space** *S*: set of all Hamiltonian cycles in given graph *G*.
- Solution set: same as S
- ▶ Neighborhood relation  $\mathcal{N}(s)$ : 2-edge-exchange
- Initialization: an Hamiltonian cycle uniformly at random.
- ▶ Step function: implemented as 2-stage process:
  - 1. select neighbor  $s' \in N(s)$  uniformly at random;
  - 2. accept as new search position with probability:

$$p(T, s, s') := egin{cases} 1 & ext{if } f(s') \leq f(s) \ \exp rac{-(f(s') - f(s))}{ au} & ext{otherwise} \end{cases}$$

(Metropolis condition), where *temperature* parameter T controls likelihood of accepting worsening steps.

**Termination:** upon exceeding given bound on run-time.

### **Great Deluge**

Another version of probabilistic iterative improvement:

Key Idea: start with a minimum water level,

- at each step a candidate move is generated at random, the move is accepted if its quality is greater than the water level.
- After the number of neighbors have been sampled at a certain water level, the water level is updated.
- The algorithm stops if we have reached the maximum water level or if we have done a certain number of non-improving solutions.

#### In greatdeluge.hh

```
InitializeRun()
{
level = initial level * this->current state cost;
}
SelectMove()
{
 this->SelectRandomMove();
}
StopCriterion()
{
 return level < min level * this->best state cost;
}
UpdateIterationCounter()
{
 if (this->number of iterations % neighbors sampled == 0)
  {
 level *= level rate;
}
AcceptableMove()
{
 return LessOrEqualThan(this->current move cost,(CFtype)0) || LessOrEqualThan((double)(
       this->current move cost + this->current state cost), level);
}
```

#### 1. Trajectory Based Metaheuristics Stochastic Local Search Simulated Annealing

Iterated Local Search Tabu Search Variable Neighborhood Search Guided Local Search

#### Inspired by statistical mechanics in matter physics:

- candidate solutions  $\cong$  states of physical system
- evaluation function  $\cong$  thermodynamic energy
- globally optimal solutions  $\cong$  ground states
- parameter  $T \cong$  physical temperature

*Note:* In physical process (*e.g.*, annealing of metals), perfect ground states are achieved by very slow lowering of temperature.





# Simulated Annealing

**Key idea:** Vary temperature parameter, *i.e.*, probability of accepting worsening moves, in Probabilistic Iterative Improvement according to annealing schedule (aka *cooling schedule*).

### Simulated Annealing (SA):

determine initial candidate solution sset initial temperature T according to annealing schedule **while** termination condition is not satisfied: **do** 

update T according to annealing schedule

- 2-stage step function based on
  - proposal mechanism (often uniform random choice from N(s))
  - acceptance criterion (often *Metropolis condition*)
- Annealing schedule

(function mapping run-time t onto temperature T(t)):

- initial temperature T<sub>0</sub> (may depend on properties of given problem instance)
- temperature update scheme (e.g., linear cooling: T<sub>i+1</sub> = T<sub>0</sub>(1 - i/I<sub>max</sub>),

geometric cooling:  $T_{i+1} = \alpha \cdot T_i$ )

- number of search steps to be performed at each temperature (often multiple of neighborhood size)
- may be static or dynamic
- seek to balance moderate execution time with asymptotic behavior properties
- Termination predicate: often based on acceptance ratio, i.e., ratio accepted / proposed steps or number of idle iterations

#### Example: Simulated Annealing for TSP

Extension of previous PII algorithm for the TSP, with

- proposal mechanism: uniform random choice from 2-exchange neighborhood;
- ► acceptance criterion: Metropolis condition (always accept improving steps, accept worsening steps with probability exp [-(f(s') - f(s))/T]);
- ▶ annealing schedule: geometric cooling T := 0.95 · T with n · (n − 1) steps at each temperature (n = number of vertices in given graph), T<sub>0</sub> chosen such that 97% of proposed steps are accepted;
- termination: when for five successive temperature values no improvement in solution quality and acceptance ratio < 2%.</p>

Improvements:

- neighborhood pruning (e.g., candidate lists for TSP)
- ▶ greedy initialization (*e.g.*, by using NNH for the TSP)
- low temperature starts (to prevent good initial candidate solutions from being too easily destroyed by worsening steps)

# Profiling



#### 1. Trajectory Based Metaheuristics

Stochastic Local Search Simulated Annealing

### Iterated Local Search

Tabu Search Variable Neighborhood Search Guided Local Search

### Iterated Local Search

Key Idea: Use two types of LS steps:

- subsidiary local search steps for reaching local optima as efficiently as possible (intensification)
- perturbation steps for effectively escaping from local optima (diversification).

*Also:* Use acceptance criterion to control diversification *vs* intensification behavior.

```
Iterated Local Search (ILS):
determine initial candidate solution s
perform subsidiary local search on s
while termination criterion is not satisfied do
r := s
perform perturbation on s
perform subsidiary local search on s
based on acceptance criterion,
keep s or revert to s := r
```

#### Note:

- Subsidiary local search results in a local minimum.
- ► ILS trajectories can be seen as walks in the space of local minima of the given evaluation function.
- Perturbation phase and acceptance criterion may use aspects of search history (i.e., limited memory).
- In a high-performance ILS algorithm, subsidiary local search, perturbation mechanism and acceptance criterion need to complement each other well.

#### Subsidiary local search:

- More effective subsidiary local search procedures lead to better ILS performance. Example: 2-opt vs 3-opt vs LK for TSP.
- Often, subsidiary local search = iterative improvement, but more sophisticated LS methods can be used. (*e.g.*, Tabu Search).

#### Perturbation mechanism:

- Needs to be chosen such that its effect *cannot* be easily undone by subsequent local search phase.
   (Often achieved by search steps larger neighborhood.)
   *Example:* local search = 3-opt, perturbation = 4-exchange steps in ILS for TSP.
- A perturbation phase may consist of one or more perturbation steps.
- ▶ Weak perturbation ⇒ short subsequent local search phase; but: risk of revisiting current local minimum.
- ► Strong perturbation ⇒ more effective escape from local minima; but: may have similar drawbacks as random restart.
- Advanced ILS algorithms may change nature and/or strength of perturbation adaptively during search.

### Components

#### Acceptance criteria:

Always accept the best of the two candidate solutions

 $\Rightarrow$  ILS performs Iterative Improvement in the space of local optima reached by subsidiary local search.

- Always accept the most recent of the two candidate solutions
   ILS performs random walk in the space of local optima reached by subsidiary local search.
- Intermediate behavior: select between the two candidate solutions based on the *Metropolis criterion* (*e.g.*, used in *Large Step Markov Chains* [Martin et al., 1991].
- Advanced acceptance criteria take into account search history, e.g., by occasionally reverting to *incumbent solution*.

### Examples

Example: Iterated Local Search for the TSP (1)

- **Given:** TSP instance  $\pi$ .
- **Search space:** Hamiltonian cycles in  $\pi$ .
- **Subsidiary local search:** Lin-Kernighan variable depth search algorithm
- Perturbation mechanism: 'double-bridge move' = particular 4-exchange step:



 Acceptance criterion: Always return the best of the two given candidate round trips.

#### 1. Trajectory Based Metaheuristics

Stochastic Local Search Simulated Annealing Iterated Local Search

#### Tabu Search

Variable Neighborhood Search Guided Local Search

**Key idea:** Avoid repeating history (memory) How can we remember the history without

- memorizing full solutions (space)
- computing hash functions (time)

 $\rightsquigarrow$  use attributes

### Tabu Search

Key idea: Use aspects of search history (memory) to escape from local minima.

- Associate tabu attributes with candidate solutions or solution components.
- Forbid steps to search positions recently visited by underlying iterative best improvement procedure based on tabu attributes.

#### Tabu Search (TS):

```
determine initial candidate solution s
While termination criterion is not satisfied:
determine set N' of non-tabu neighbors of s
choose a best candidate solution s' in N'
update tabu attributes based on s'
s := s'
```

Example: Tabu Search for SAT

- **Search space:** set of all complete assignments of *X*.
- **Solution set:** models of the formula.
- Neighborhood relation: 1-flip
- ▶ **Memory:** Associate tabu status (Boolean value) with each pair (literal,value) (*x*, *val*).
- Initialization: a random assignment
- Search steps:
  - pairs (x, v) are tabu if they have been changed in the last tt steps;
  - neighboring assignments are admissible if they can be reached by changing a non-tabu pair or have fewer unsatisfied constraints than the best assignments seen so far (*aspiration criterion*);
  - choose uniformly at random admissible neighbors with minimal number of unsatisfied constraints.
- Termination: upon finding a feasible assignment or after given bound on number of search steps has been reached or after a number of idle iterations

#### Note:

- Admissible neighbors of s: Non-tabu search positions in N(s)
- ► Tabu tenure: a fixed number of subsequent search steps for which the last search position or the solution components just added/removed from it are declared tabu
- Aspiration criterion (often used): specifies conditions under which tabu status may be overridden (*e.g.*, if considered step leads to improvement in incumbent solution).
- Crucial for efficient implementation:
  - efficient best improvement local search

     pruning, delta updates, (auxiliary) data structures
  - efficient determination of tabu status: store for each variable x the number of the search step when its value was last changed *it*<sub>x</sub>; x is tabu if *it* - *it*<sub>x</sub> < *tt*, where *it* = current search step number.

#### 1. Trajectory Based Metaheuristics

Stochastic Local Search Simulated Annealing Iterated Local Search Tabu Search

#### Variable Neighborhood Search

Guided Local Search

Variable Neighborhood Search is a method based on the systematic change of the neighborhood during the search.

#### Central observations

- a local minimum w.r.t. one neighborhood function is not necessarily locally minimal w.r.t. another neighborhood function
- ▶ a global optimum is locally optimal w.r.t. all neighborhood functions

#### Key principle: change the neighborhood during the search

- Several adaptations of this central principle
  - (Basic) Variable Neighborhood Descent (VND)
  - Variable Neighborhood Search (VNS)
  - Reduced Variable Neighborhood Search (RVNS)
  - Variable Neighborhood Decomposition Search (VNDS)
  - Skewed Variable Neighborhood Search (SVNS)
- Notation
  - ▶  $\mathcal{N}_k$ ,  $k = 1, 2, ..., k_m$  is a set of neighborhood functions
  - $N_k(s)$  is the set of solutions in the *k*-th neighborhood of *s*

How to generate the various neighborhood functions?

- for many problems different neighborhood functions (local searches) exist / are in use
- change parameters of existing local search algorithms
- ▶ use *k*-exchange neighborhoods; these can be naturally extended
- many neighborhood functions are associated with distance measures; in this case increase the distance

### Basic Variable Neighborhood Descent

### **Procedure** BVND input : $\mathcal{N}_k$ , $k = 1, 2, ..., k_{max}$ , and an initial solution *s* output: a local optimum *s* for $\mathcal{N}_k$ , $k = 1, 2, ..., k_{max}$ $k \leftarrow 1$

repeat

```
| \begin{array}{c} s' \leftarrow \mathsf{FindBestNeighbor}(s, \mathcal{N}_k) \\ \mathbf{if} \ f(s') < f(s) \ \mathbf{then} \\ \  \  \left\lfloor \begin{array}{c} s \leftarrow s' \\ (k \leftarrow 1) \end{array} \right. \\ \mathbf{else} \\ \  \  \left\lfloor \begin{array}{c} k \leftarrow k+1 \end{array} \right. \\ \mathbf{until} \ k = k_{max}, \end{array} \right.
```

# Variable Neighborhood Descent

**Procedure VND** input :  $\mathcal{N}_k$ ,  $k = 1, 2, ..., k_{max}$ , and an initial solution *s* output: a local optimum *s* for  $\mathcal{N}_k$ ,  $k = 1, 2, ..., k_{max}$  $k \leftarrow 1$ 

repeat

```
s' \leftarrow lterativeImprovement(s, \mathcal{N}_k)
if f(s') < f(s) then
s \leftarrow s'
k \leftarrow 1
else
k \leftarrow k + 1
until k = k_{max}:
```

- ► Final solution is locally optimal w.r.t. all neighborhoods
- First improvement may be applied instead of best improvement
- ► Typically, order neighborhoods from smallest to largest
- ► If iterative improvement algorithms II<sub>k</sub>, k = 1,..., k<sub>max</sub> are available as black-box procedures:
  - order black-boxes
  - apply them in the given order
  - possibly iterate starting from the first one
  - order chosen by: solution quality and speed

### Basic Variable Neighborhood Search

Procedure BVNS **input** :  $\mathcal{N}_k$ ,  $k = 1, 2, \dots, k_{max}$ , and an initial solution s **output**: a local optimum s for  $\mathcal{N}_k$ ,  $k = 1, 2, \ldots, k_{max}$ repeat  $k \leftarrow 1$ repeat  $| s' \leftarrow \mathsf{RandomPicking}(s, \mathcal{N}_k)$  $s'' \leftarrow \text{IterativeImprovement}(s', \mathcal{N}_k)$ if f(s'') < f(s) then  $\left[\begin{array}{c} \dot{s} \leftarrow s''\\ k \leftarrow 1 \end{array}\right]$ else until Termination Condition;

To decide:

- which neighborhoods
- how many
- which order
- which change strategy

► Extended version: parameters k<sub>min</sub> and k<sub>step</sub>; set k ← k<sub>min</sub> and increase by k<sub>step</sub> if no better solution is found (achieves diversification)

#### 1. Trajectory Based Metaheuristics

Stochastic Local Search Simulated Annealing Iterated Local Search Tabu Search Variable Neighborhood Search Guided Local Search

## **Guided Local Search**

- ► **Key Idea:** Modify the evaluation function whenever a local optimum is encountered.
- Associate weights (penalties) with solution components; these determine impact of components on evaluation function value.
- Perform Iterative Improvement; when in local minimum, increase penalties of some solution components until improving steps become available.

#### Guided Local Search (GLS):

```
determine initial candidate solution s
initialize penalties
while termination criterion is not satisfied do
compute modified evaluation function g' from g
based on penalties
perform subsidiary local search on s
using evaluation function g'
update penalties based on s
```

#### Guided Local Search (continued)

Modified evaluation function:

$$g'(s) := f(s) + \sum_{i \in SC(s)} \text{penalty}(i),$$

where SC(s) is the set of solution components used in candidate solution *s*.

- Penalty initialization: For all i: penalty(i) := 0.
- Penalty update in local minimum s: Typically involves penalty increase of some or all solution components of s; often also occasional penalty decrease or penalty smoothing.
- **Subsidiary local search:** Often *Iterative Improvement*.

#### Potential problem:

Solution components required for (optimal) solution may also be present in many local minima.

#### Possible solutions:

- A: Occasional decreases/smoothing of penalties.
- **B:** Only increase penalties of solution components that are least likely to occur in (optimal) solutions.

#### Implementation of **B**:

Only increase penalties of solution components *i* with maximal utility [Voudouris and Tsang, 1995]:

$$\mathtt{util}(s,i) := \frac{f_i(s)}{1 + \mathtt{penalty}(i)}$$

where  $f_i(s)$  is the solution quality contribution of *i* in *s*.

#### Example: Guided Local Search (GLS) for the TSP

[Voudouris and Tsang 1995; 1999]

- Given: TSP instance  $\pi$
- **Search space:** Hamiltonian cycles in  $\pi$  with *n* vertices;
- Neighborhood: 2-edge-exchange;
- Solution components edges of π; f<sub>e</sub>(G, p) := w(e);
- Penalty initialization: Set all edge penalties to zero.
- Subsidiary local search: Iterative First Improvement.
- ► **Penalty update:** Increment penalties of all edges with maximal utility by  $\lambda := 0.3 \cdot \frac{w(s_{2-opt})}{n}$

where  $s_{2-opt} = 2$ -optimal tour.

### Summary

#### 1. Trajectory Based Metaheuristics

Stochastic Local Search Simulated Annealing Iterated Local Search Tabu Search Variable Neighborhood Search Guided Local Search

1. Trajectory Based Metaheuristics

Stochastic Local Search Simulated Annealing Iterated Local Search Tabu Search Variable Neighborhood Search Guided Local Search

#### 2. Population Based Metaheuristics

Evolutionary Algorithms Ant Colony Optimization

1. Trajectory Based Metaheuristics Stochastic Local Search Simulated Annealing

> Tabu Search Variable Neighborhood Search Guided Local Search

1. Trajectory Based Metaheuristics Stochastic Local Search

> Simulated Annealing Iterated Local Search Tabu Search Variable Neighborhood Search Guided Local Search