DM841 Discrete Optimization

Part I

Lecture 5 More Examples

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Outline

1. Examples

n-Queens, Grocery, Magic Square

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1. Examples

n-Queens, Grocery, Magic Square

Constraints: No Attack

not in same column

- by choice of variables
- not in same row
 - $x_i \neq x_j$ for $i \neq j$
- not in same diagonal

•
$$x_i - i \neq x_j - j$$
 for $i \neq j$
• $x_i - j \neq x_i - i$ for $i \neq j$

• $3 \cdot n \cdot (n-1)$ constraints

Fewer Constraints...

Sufficient by symmetry
 i < *j* instead of *i* ≠ *j* Constraints
 x_i ≠ *x_j for i* < *j j i* < *j i* < *j j*

$$x_i - i \neq x_j - j for i < j x_i - j \neq x_j - i for i < j$$

•
$$3/2 \cdot n \cdot (n-1)$$
 constraints

Even Fewer Constraints

Not same row constraint

 $x_i \neq x_j$ for i < jmeans: values for variables pairwise distinct

Constraints

- $x_i i \neq x_j j$ for i < j
- $x_i j \neq x_j i$ for i < j

Pushing it Further...

 Yes, also diagonal constraints can be captured by distinct constraints

see assignment

distinct(x0, x1, ..., x7) distinct(x0-0, x1-1, ..., x7-7) distinct(x0+0, x1+1, ..., x7+7) Script: Variables

Queens(void) : q(*this,8,0,7) {

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...

}

Script: Constraints

```
Queens(void) : q(*this,8,0,7) {
    distinct(*this, q);
    for (int i=0; i<8; i++)
        for (int j=i+1; j<8; j++) {
        rel post(*this, x[i]-i != x[j]-j);
        post(*this, x[i]-j != x[j]-i);
        }
    ...
}</pre>
```

Script: Branching

```
Queens(void) : q(*this,8,0,7) {
    ...
    branch(*this, q,
        INT_VAR_NONE,
        INT_VAL_MIN);
}
```

Good Branching?

Naïve is not a good strategy for branching

Try the following (see assignment)

- first fail
- place queen as much in the middle of a row
- place queen in knight move fashion

Summary 8 Queens

Variables

- model should require few variables
- good: already impose constraints

Constraints

- do not post same constraint twice
- try to find "big" constraints subsuming many small constraints
 - more efficient
 - often, more propagation (to be discussed)



Grocery

- Kid goes to store and buys four items
- Cashier: that makes \$7.11
- Kid: pays, about to leave store
- Cashier: hold on, I multiplied!
 - let me add!

wow, sum is also \$7.11

You: prices of the four items?

Model

Variables

- for each item A, B, C, D
- take values between {0, ..., 711}
- compute with cents: allows integers
- Constraints
 - A + B + C + D = 711

The unique solution (upon the symmetry breaking of slide 87) is: A=120, B=125, C=150, D=316.

Γ.

```
Script
```

```
class Grocery : public Space {
protected:
    IntVarArray abcd;
    const int s = 711;
    const int p = s * 100 * 100 * 100;
public:
    Grocery(void) ... { ... }
    •••
}
```

Script: Variables

Grocery(void) : abcd(*this,4,0,711) {

•••

}

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Script: Sum

```
…
// Sum of all variables is s
linear(this, abcd, IRT_EQ, s);
```

Script: Product

IntVar t1(*this,1,p); IntVar t2(*this,1,p); IntVar t3(*this,p,p);

Branching

- Bad idea: try values one by one
- Good idea: split variables
 - for variable x
 - with $m = (\min(x) + \max(x)) / 2$
 - branch x < m or $x \ge m$
- Typically good for problems involving arithmetic constraints
 - exact reason needs to be explained later

Script: Branching

Search Tree

- 2829 nodes for first solution
- Pretty bad...

Better Heuristic?

 Try branches in different order split with larger interval first
 try: INT_VAL_SPLIT_MAX
 Search tree: 2999 nodes

worse in this case

Symmetries

- Interested in values for A, B, C, D
- Model admits equivalent solutions
 - interchange values for A, B, C, D
- We can add order A, B, C, D: A ≤ B ≤ C ≤ D
- Called "symmetry breaking constraint"

Script: Symmetry Breaking

•••

rel(this, a, IRT_LQ, b);
rel(this, b, IRT_LQ, c);
rel(this, c, IRT_LQ, d);

•••

Effect of Symmetry Breaking

Search tree size 308 nodes

Let us try INT_VAL_SPLIT_MAX again

- tree size 79 nodes!
- interaction between branching and symmetry breaking
- other possibility: $A \ge B \ge C \ge D$
- we need to investigate more (later)!

Any More Symmetries?

Observe: 711 has prime factor 79

that is: 711 = 79 × 9

Assume: A can be divided by 79

add: A = 79 × X

for some finite domain var X

- remove A ≤ B
- the remaining B, C, D of course can still be ordered

Any More Symmetries?

In Gecode

IntVar x(*this,1,p);

IntVar sn(*this,79,79);

mult(*this, x, sn, a);

Search tree 44 nodes!

now we are talking!

Summary: Grocery

Branching: consider also

- how to partition domain
- in which order to try alternatives
- Symmetry breaking
 - can reduce search space
 - might interact with branching
 - typical: order variables in solutions
- Try to really understand problem!

Domination Constraints

- In symmetry breaking, prune solutions without interest
- Similarly for best solution search
 - typically, interested in just one best solution
 - impose constraints to prune some solutions with same "cost"

Another Observation

• Multiplication decomposed as $A \cdot B = T_1$ $C \cdot D = T_2$ $T_1 \cdot T_2 = P$

What if

- $A \cdot B = T_1 \quad T_1 \cdot C = T_2 \quad T_2 \cdot D = P$
- propagation changes: 355 nodes
- propagation is not compositional!
- another point to investigate



2	9	4
7	5	3
6	1	8

Unique solution for n=3, upon the symmetry breaking of slide 99.

Magic Squares

Find an *n*×*n* matrix such that

- every field is integer between 1 and n^2
- fields pairwise distinct
- sums of rows, columns, two main diagonals are equal
- Very hard problem for large n
- Here: we just consider the case n=3

Model

For each matrix field have variable x_{ii}

• $x_{ij} \in \{1, ..., 9\}$

One additional variable s for sum

■ *s* ∈ {1, .., 9×9}

- All fields pairwise distinct
 - distinct(x_{ij})
- For each row i have constraint

• $x_{i0} + x_{i1} + x_{i2} = s$

columns and diagonals similar

Script

- Straightforward
- Branching strategy
 - first-fail
 - split again: arithmetic constraints
 - try to come up with something that is really good!

Generalize it to arbitrary n

Symmetries

- Clearly, we can require for first row that first and last variable must be in order
- Also, for opposing corners
- In all (other combinations possible)
 - $x_{00} < x_{02}$
 - $x_{02} < x_{20}$
 - $x_{00} < x_{22}$

Important Observation

We know the sum of all fields 1 + 2 + ... + 9 = 9(9+1)/2=45
We "know" the sum of one row *s*We know that we have three rows 3×s = 45

Implied Constraints

The constraint model already implies

3×*s* = 45

implies solutions are the same

- However, adding a propagator for the constraint drastically improves propagation
- Often also: redundant or implied constraint

Effect

- Simple model
- Symmetry breaking
- Implied constraint

92 nodes 29 nodes 6 nodes

Summary: Magic Squares

Add implied constraints

- are implied by model
- increase constraint propagation
- reduce search space
- require problem understanding
- Also as usual
 - break symmetries
 - choose appropriate branching

Outlook...

Common modeling principles

- what are the variables
- finding the constraints
- finding the propagators
- implied (redundant) constraints
- finding the branching
- symmetry breaking

Modeling Strategy

Understand problem

- identify variables
- identify constraints
- identify optimality criterion
- Attempt initial model simple?
 - try on examples to assess correctness
- Improve model

- much harder!
- scale up to real problem size

Summary

Experiment with:

- different branching strategies
- different models (eg, adding redundat constraints)
- different propagation strength (to come)