DM204, 2010 SCHEDULING, TIMETABLING AND ROUTING

Local Search Methods for VRPTW

Marco Chiarandini

Department of Mathematics & Computer Science University of Southern Denmark

Course Overview

- ✓ Problem Introduction
 - ✓ Scheduling classification
 - Scheduling complexity
 - ✓ RCPSP
- ✓ General Methods
 - ✓ Integer Programming
 - ✓ Constraint Programming
 - ✓ Heuristics
 - ✓ Dynamic Programming
 - ✔ Branch and Bound

✓ Scheduling Models

- ✓ Single Machine
- ✓ Parallel Machine and Flow Shop
- ✓ Job Shop
- ✓ Resource-Constrained Project Scheduling
- Timetabling
 - Reservations and Education
 - ✓ Course Timetabling
 - ✓ Workforce Timetabling
 - Crew Scheduling
- Vehicle Routing
 - Capacitated Models
 - ✓ Time Windows models
 - Rich Models

Outline

1. Variants of VRP

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Rich VRP

Definition

Rich Models are non idealized models that represent the application at hand in an adequate way by including all important optimization criteria, constraints and preferences [Hasle et al., 2006]

Solution

- Exact methods are often impractical:
 - instances are too large
 - decision support systems require short response times
- Metaheuristics based on local search components are mostly used

VRP with Backhauls

Further Input from CVRP:

a partition of customers:

$$L = \{1, ..., n\}$$
 Lineahaul customers (deliveries)
 $B = \{n + 1, ..., n + m\}$ Backhaul customers (collections)

• precedence constraint: in a route, customers from L must be served before customers from B

Task: Find a collection of K simple circuits with minimum costs, such that:

- each circuit visit the depot vertex
- each customer vertex is visited by exactly one circuit; and
- the sum of the demands of the vertices visited by a circuit does not exceed the vehicle capacity Q.
- in any circuit all the linehaul customers preced the backhaul customers, if any.

VRP with Pickup and Delivery

Further Input from CVRP:

- each customer i is associated with quantities d_i and p_i to be delivered and picked up, resp.
- for each customer i, Oi denotes the vertex that is the origin of the delivery demand and D_i denotes the vertex that is the destination of the pickup demand

Task:

Find a collection of K simple circuits with minimum costs, such that:

- each circuit visit the depot vertex
- each customer vertex is visited by exactly one circuit; and
- the current load of the vehicle along the circuit must be non-negative and may never exceed Q
- for each customer i, the customer O_i when different from the depot, must be served in the same circuit and before customer i
- for each customer i, the customer D_i when different from the depot, must be served in the same circuit and after customer i

Multiple Depots VRP

Further Input from CVRP:

- multiple depots to which customers can be assigned
- a fleet of vehicles at each depot

Task:

Find a collection of K simple circuits for each depot with minimum costs, such that:

- each circuit visit the depot vertex
- each customer vertex is visited by exactly one circuit; and
- the current load of the vehicle along the circuit must be non-negative and may never exceed Q
- vehicles start and return to the depots they belong

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Vertex set V = \{1, 2, ..., n\} and V_0 = \{n + 1, ..., n + m\}
Route i defined by R_i = \{l, 1, ..., l\}
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Periodic VRP

Further Input from CVRP:

planning period of M days

Task:

Find a collection of K simple circuits with minimum costs, such that:

- each circuit visit the depot vertex
- each customer vertex is visited by exactly one circuit; and
- the current load of the vehicle along the circuit must be non-negative and may never exceed Q
- A vehicle may not return to the depot in the same day it departs.
- Over the M-day period, each customer must be visited / times, where 1 < I < M.

Three phase approach:

Generate feasible alternatives for each customer.

Example, M = 3 days $\{d1, d2, d3\}$ then the possible combinations are:

0
$$\rightarrow$$
 000; 1 \rightarrow 001; 2 \rightarrow 010; 3 \rightarrow 011; 4 \rightarrow 100; 5 \rightarrow 101; 6 \rightarrow 110; 7 \rightarrow 111

Customer	Diary De- mand	Number of Visits	Number of Combina-	Possible Combina-
			tions	tions
1	30	1	3	1,2,4
2	20	2	3	3,4,6
3	20	2	3	3,4,6
4	30	2	3	1,2,4
5	10	3	1	7

- 2. Select one of the alternatives for each customer, so that the daily constraints are satisfied. Thus, select the customers to be visited in each day.
- 3. Solve the vehicle routing problem for each day.

Split Delivery VRP

Constraint Relaxation: it is allowed to serve the same customer by different vehicles. (necessary if $d_i > Q$)

Task:

Find a collection of K simple circuits with minimum costs, such that:

- each circuit visit the depot vertex
- the current load of the vehicle along the circuit must be non-negative and may never exceed Q

Note: a SDVRP can be transformed into a VRP by splitting each customer order into a number of smaller indivisible orders [Burrows 1988].

Inventory VRP

Input:

- a facility, a set of customers and a planning horizon T
- r_i product consumption rate of customer i (volume per day)
- C_i maximum local inventory of the product for customer i
- a fleet of M homogeneous vehicles with capacity Q

Task:

Find a collection of K daily circuits to run over the planing horizon with minimum costs and such that:

- each circuit visit the depot vertex
- no customer goes in stock-out during the planning horizon
- the current load of the vehicle along the circuit must be non-negative and may never exceed Q

Other VRPs

VRP with Satellite Facilities (VRPSF)

Possible use of satellite facilities to replenish vehicles during a route.

Open VRP (OVRP)

The vehicles do not need to return at the depot, hence routes are not circuits but paths

Dial-a-ride VRP (DARP)

- It generalizes the VRPTW and VRP with Pick-up and Delivery by incorporating time windows and maximum ride time constraints
- It has a human perspective
- Vehicle capacity is normally constraining in the DARP whereas it is often redundant in PDVRP applications (collection and delivery of letters and small parcels)