# DM545/DM871 – Linear and integer programming

Sheet 4, Autumn 2024

## Solution:

## Included.

Exercises with the symbol  $^+$  are to be done at home before the class. Exercises with the symbol  $^*$  will be tackled in class and should be at least read at home. The remaining exercises are left for self training after the exercise class.

## Exercise 1<sup>+</sup>

Solve the systems  $\mathbf{y}^T E_1 E_2 E_3 E_4 = [1 \ 2 \ 3]$  and  $E_1 E_2 E_3 E_4 \mathbf{d} = [1 \ 2 \ 3]^T$  with

$E_1 =$	$\begin{bmatrix} 1 & 3 & 0 \\ 0 & 0.5 & 0 \\ 0 & 4 & 1 \end{bmatrix}$	$E_2 = \begin{bmatrix} 2 & 0 & 0 \\ 1 & 1 & 0 \\ 4 & 0 & 1 \end{bmatrix}$	$E_3 = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 3 \\ 0 & 0 & 1 \end{bmatrix}$	$E_4 =$	-0.5 3 1	$\begin{bmatrix} 0 & 0 \\ 1 & 0 \\ 0 & 1 \end{bmatrix}$
						0 1

## Solution:

This exercise is to show that the two systems can be solved quite easily. Let's take first  $\mathbf{y}^T E_1 E_2 E_3 E_4 = [1 \ 2 \ 3]$ , we use the backward transformation and solve the sequence of linear systems:

$$\mathbf{u}^{T} E_{4} = [1 \ 2 \ 3], \quad \mathbf{v}^{T} E_{3} = \mathbf{u}^{T}, \quad \mathbf{w}^{T} E_{2} = \mathbf{v}^{T}, \quad \mathbf{y}^{T} E_{1} = \mathbf{w}^{T}$$
$$\mathbf{u}^{T} \begin{bmatrix} -0.5 & 0 & 0\\ 3 & 1 & 0\\ 1 & 0 & 1 \end{bmatrix} = [1, 2, 3]$$

Since the eta matrices have always one 1 in two columns then the solution can be read up easily. From the third column we find  $u_3 = 3$ . From the second column, we find  $u_2 = 2$ . Substituting in the first column, we find  $-0.5u_1 + 3 * 2 + 1 * 3 = 1$ , which yields  $u_1 = 16$ . The next system is:

$$\mathbf{v}^{T} \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 3 \\ 0 & 0 & 1 \end{bmatrix} = [16, 2, 3]$$

From the first column we get  $v_1 = 16$ , from the second column  $v_2 = 2$  from the last column  $v_3 = -19$ . The next:

$$\mathbf{w}^{T} \begin{bmatrix} 2 & 0 & 0 \\ 1 & 1 & 0 \\ 4 & 0 & 1 \end{bmatrix} = [16, 2, -19]$$

which gives  $\mathbf{w} = [45, 2, -19]$ . The next:

$$\mathbf{y}^{T} \begin{bmatrix} 1 & 3 & 0 \\ 0 & 0.5 & 0 \\ 0 & 4 & 1 \end{bmatrix} = [45, 2, -19]$$

from we finally find d = [119/2, 2, -19].

We can carry out the operations in Python for cross checking the results but the function linalg.solve would not exploit the advantages of the eta matrices when solving the system of linear equations.

```
import numpy as np
```

```
E4=np.array([[-0.5,0,0],[3,1,0],[1,0,1]])
E3=np.array([[1,0,1],[0,1,3],[0,0,1]])
E2=np.array([[2,0,0],[1,1,0],[4,0,1]])
E1=np.array([[1,3,0],[0,0.5,0],[0,4,1]])
u=np.linalg.solve(E4.T,np.array([1,2,3]))
print(u)
v=np.linalg.solve(E3.T,u)
print(v)
w=np.linalg.solve(E2.T,v)
print(w)
d=np.linalg.solve(E2.T,w)
print(d)
```

[16. 2. 3.] [ 16. 2. -19.] [ 45. 2. -19.] [ 59.5 2. -19. ]

## Exercise 2\* Sensitivity Analysis and Revised Simplex

A furniture-manufacturing company can produce four types of products using three resources.

- A bookcase requires three hours of work, one unit of metal, and four units of wood and it brings in a net profit of 19 Euro.
- A desk requires two hours of work, one unit of metal and three units of wood, and it brings in a net profit of 13 Euro.
- A chair requires one hour of work, one unit of metal and three units of wood and it brings in a net profit of 12 Euro.
- A bedframe requires two hours of work, one unit of metal, and four units of wood and it brings in a net profit of 17 Euro.
- Only 225 hours of labor, 117 units of metal and 420 units of wood are available per day.

In order to decide how much to make of each product so as to maximize the total profit, the managers solve an LP problem.

1) Write the mathematical programming formulation of the problem.

Solution:

 $\max \begin{array}{l} 19x_1 + 13x_2 + 12x_3 + 17x_4 \\ 3x_1 + 2x_2 + x_3 + 2x_4 \le 225 \\ x_1 + x_2 + x_3 + x_4 \le 117 \\ 4x_1 + 3x_2 + 3x_3 + 4x_4 \le 420 \\ x_1, x_2, x_3, x_4 \ge 0 \end{array}$ 

2) With the help of a computational environment such as Python for carrying out linear algebra operations, write the optimal tableau, which has  $x_1, x_3$  and  $x_4$  in basis.

Do this task:

- first, using the original simplex method
- second, using the revised simplex. In this case, start by writing  $A_B$ ,  $A_N$ , then calculate  $A_B^{-1}A_N$ , and the finally derive the optimal tableau and verify that the solution is indeed optimal.

## Solution:

The initial tableau is:

Ì	x_1	I	x_2	I	x_3	3	c_4	I	x_5	I	x_6	I	x_7	I	-z		b
Ì	3	I	2	I	1	I	2	I	1	I	0	I	0	I	0	22	25
					1 3	· · ·								÷.,			
1					12					1		1		1			

We know that there will be 3 variables in basis. The text of the problem tells us which these 3 variables are: 1, 3, 4. Hence,

	3	1	2		2	1	0	0	
$A_B =$	1	1	1	$A_N =$	1	0	1	0	
	4	3	4	$A_N =$	3	0	0	1	

We can calculate  $A_B^{-1}A_N$  in Python or in R:

```
> B=matrix(c(3,1,2,1,1,1,4,3,4),byrow=TRUE,ncol=3)
> B1=solve(B)
> B%*%B1 # check to make sure it is correct!
    [,1] [,2] [,3]
[1,]
      1 0 0
[2,]
       0
            1
                0
          0
[3,]
      0
                 1
> N=matrix(c( 2, 1, 0, 0, 1, 0, 1, 0, 3, 0, 0, 1),ncol=4,byrow=TRUE)
> B1%*%N
    [,1] [,2] [,3] [,4]
[1,]
      1
          1
                 2
                    -1
[2,]
      1
            0
                 4
                     -1
[3,]
     -1 -1
                -5
                      2
> cN=c(13,0,0,0)
> cB=c(19,12,17)
> cN-cB%*%B1%*%N
    [,1] [,2] [,3] [,4]
[1,]
     -1 -2 -1 -3
> cB%*%B1%*%c(225,117,420)
> 1827
```

This code gives us:

$$\bar{A} = A_B^{-1} A_N = \begin{bmatrix} 1 & 1 & 2 & -1 \\ 1 & 0 & 4 & -1 \\ -1 & -1 & -5 & 2 \end{bmatrix} \qquad x_B^* = \begin{bmatrix} x_1 \\ x_3 \\ x_4 \end{bmatrix} = A_B^{-1} b = \begin{bmatrix} 39 \\ 48 \\ 30 \end{bmatrix}$$
$$\bar{c}_N = \begin{bmatrix} \bar{c}_2 & \bar{c}_5 & \bar{c}_6 & \bar{c}_7 \end{bmatrix} = \begin{bmatrix} -1 & -2 & -1 & -3 \end{bmatrix}$$

and we can write the final tableau as:

	+-		+		+-		+-		+-		+-		+		+	
x_1	I	x_2	I	x_3	L	x_4	L	x_5		x_6	I	x_7	I	-z	I	Ъ
																 39
0	I.	1	I	1	L	0	L	0		4		-1	T	0		48
0	T	-1	T	0	L	1		-1		-5		2	T	0	I	30
0	T	-1	T	0	L	0	L	-2		-1		-3	T	1	T	-1827
	+-		+		+-		+-		+-		+-		+		+	

Since all reduced costs are negative then the tableau and the corresponding solution are optimal.

3) What is the increase in price (reduced cost) that would make product  $x_2$  worth to be produced?

## Solution:

The increase in price of a quantity strictly larger than 1 would make the product 2 worth being produced. Indeed, let  $c'_2 = c_2 + \delta$  be the new price. We know that the coefficient in the objective function goes in the reduced cost calculation multiplied by 1. Hence, to have a positive reduced cost we have:

 $-1 + \delta > 0 \implies \delta > 1$ 

We could also recalcuate the reduced cost from scratch using the multipliers  $\pi$ :  $c'_2 + \sum_{i=1}^{3} \pi_i a_{i2}$ . The value of  $\pi_i$  are read from the final tableau and they correspond to the reduced costs of the slack variables, ie, (-2, -1, -3).

4) What is the marginal value (shadow price) of an extra hour of work or amount of metal and wood?

## Solution:

The marginal values are the values of the dual variable  $y_1, y_2, y_3$ . From the strong duality theorem, we know that  $y_i = -\pi_i = -\bar{c}_{n+i}$ , i = 1..m. Hence,  $\mathbf{y} = (2, 1, 3)$ .

An extra hour of work has marginal value of 2, that is, having one unit more of work would improve the revenue by 2. For the other two resources the marginal values are 1 and 3, respectively.

We can cross check these conclusions: by the complementary slackness theorem, the fact that all three dual variables are strictly positive indicates that all three constraints in the primal are active $\equiv$ tight $\equiv$  binding. Hence, it makes sense to have that an increase in the capacity of those constraints implies an increase in the profit. The conclusion that all three constraints are tight can be also reached by the fact that the slack variables are 0 in the final tableau. If some constraint was not tight, then the marginal value of the corresponding resource would be zero since an increase in its capacity does not imply an immediate improvement in total profit.

5) Are all resources totally utilized, i.e. are all constraints "binding", or is there slack capacity in some of them? Answer this question in the light of the complementary slackness theorem.

## Solution:

Since all dual variables are strictly larger than zero, then all constraints are binding. Indeed for the complementary slackness theorem, we have that:

$$\left(b_i - \sum_{j=1}^n a_{ij} x_j^*\right) y_i^* = 0, \quad i = 1, \dots, m$$

6) From the economical interpretation of the dual why product *x*<sub>2</sub> is not worth producing? What is its imputed cost?

## Solution:

It is not worth producing 2 because  $\sum_i y_i a_{i2} > c_2$ , that is, we are better off selling the raw materials to produce the product. Indeed  $y_i$  is the price of one unit of resource i and  $a_{i2}$  is the amount of i necessary to produce 2.

$$\sum_{i} y_{i} a_{i2} = 2 * (2) + 1 * (1) + 3 * (3) = 14 > 13$$

Perform a sensitivity analysis for the following variants:

7) The net profit brought in by each desk increases from 13 Euro to 15 Euro.

## Solution:

We saw earlier that if the price of product 2 increases by more than 1 then the reduced cost becomes positive and it enters the basis. We can iterate the revised simplex as follows:

Step 1 and 2 to determine the entering varible are already done in the point a) above.

We need to do Step 3 to determine the leaving variable: we need to find the constraint that limit the increase of  $x_2$ ,  $\theta$ . We solve first  $A_B d = a$  in d. Here, **a** is the column of the matrix A (augmented with the slack variables) from the initial tableau corresponding to the entering variable  $x_2$ . We use the inverse of  $A_B$  calculated earlier in a) above in R:

> B1%\*%c(2,1,3)
 [,1]
[1,] 1
[2,] 1
[3,] -1

that is

$$\mathbf{d} = A_B^{-1} \mathbf{a} = A_B^{-1} \begin{bmatrix} 2\\1\\3 \end{bmatrix} = \begin{bmatrix} 1\\1\\-1 \end{bmatrix}$$

Then the new solutioon  $x_B$  is derived from the old one by means of **d** and the increase  $\theta$ :

$$x_B = \begin{bmatrix} x_1 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 39 \\ 48 \\ 30 \end{bmatrix} - \begin{bmatrix} 1 \\ 1 \\ -1 \end{bmatrix} \theta \ge 0$$

The increase  $\theta$  must be such that the value of the variables still remains feasible, ie,  $x_i \ge 0$ . Hence  $\theta \le 39$  and the leaving variable is  $x_1$ , since it is the one that goes to zero. The new solutions is

$$x_{B} = \begin{bmatrix} x_{2} \\ x_{3} \\ x_{4} \end{bmatrix} = \begin{bmatrix} 39 \\ 48 \\ 30 \end{bmatrix} - \begin{bmatrix} 1 \\ 1 \\ -1 \end{bmatrix} \theta = \begin{bmatrix} 39 - 39 \\ 48 - 39 \\ 30 + 39 \end{bmatrix} = \begin{bmatrix} 0 \\ 9 \\ 69 \end{bmatrix}$$

and the objective value:

```
> c=c(19,15,12,17)
> c%*%c(0,39,9,69)
      [,1]
[1,] 1866
```

8) The availability of metal increases from 117 to 125 units per day

## Solution:

This is a change in the RHS term of constraint 2. The optimality of the current solution does not change, since all reduced costs stay negative, but we need to check if it is still feasible. We need to look at the final tableau and recompute the *b* terms of all constraints. We can do this with  $A_B^{-1}b$ :

```
> b=c(225,125,420)
> B1%*%b
       [,1]
[1,] 55
[2,] 80
[3,] -10
```

The last cosntraint becomes negative, hence we need to iterate with the dual simplex.

9) The company may also produce coffee tables, each of which requires three hours of work, one unit of metal, two units of wood and bring in a net profit of 14 Euro.

Solution:

We need to check if the reduced cost of the new variable would become positive by computing  $c_0 + \sum_i \pi_i a_{ij}$ :

> 14-3\*2-1\*1-2\*3 [1] 1

which is positive, hence we need to iterate as done in point 1).

10) The number of chairs produced must be at most five times the numbers of desks

#### Solution:

This corresponds to introduce a new constraint:  $x_3 \le 5x_2$ . In the new standard form we have a new slack variable  $x_8$ . Adding the constraint in the tableau and bringing back the tableau in canonical standard form we observe that a RHS term becomes negative. Hence, we need to iterate with the dual simplex. After on iteration with the dual simplex, the final tableau becomes:

1	0	0	0	1	4/3	-5/6	1/6	0	31
0	1	0	0	0	2/3	-1/6	-1/6	0	8
0	0	0	1	-1	-13/3	11/6	-1/6	0	38
					10/3				
0	0	0	0	-2	-1/3	-19/6	-1/6	1	-1819

If after the introduction of the constraint the current solution stayed feasible then, we would have needed to check whether it was also optimal. We can either repeat the steps done at part 1 above to compute the new reduced costs or we can include the new row in the final tableau and proceed to put the tableau in canonical form. Then we look at the value of the reduced costs.

## **Exercise 3** Factory Planning and Machine Maintenance

Tasks 1-3 of Factory Planning and Maintainance Case. See Unit 7.