



# DM537

## Object-Oriented Programming

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# **STATIC FUNCTIONS FOR RECURSIVE DATA STRUCTURES**

# List ADT: Implementation 5

- Implementation 5:

```
public class RecursiveList<E> implements List<E> {  
    private ListNode<E> head = null;  
    public E get(int i) {  
        if (i < 0) { throw new IllegalArgumentException(); }  
        return get(this.head, i);  
    }  
    public static <E> E get(ListNode<E> node, int i) {  
        if (node == null) { throw new Index...Exception(); }  
        if (i == 0) { return node.getElem(); }  
        return get(node.getNext(), i-1);  
    }  
    ...  
}
```

# List ADT: Implementation 5

- Implementation 5 (continued):

```
public class RecursiveList<E> implements List<E> { ...
    public void set(int i, E elem) {
        if (i < 0) { throw new IllegalArgumentException(); }
        set(this.head, i, elem);
    }
    public static <E> void set(ListNode<E> node, int i, E elem) {
        if (node == null) { throw new Index...Exception(); }
        if (i == 0) { node.setElem(elem); }
        else { set(node.getNext(), i-1, elem); }
    }
    ...
}
```

# List ADT: Implementation 5

- Implementation 5 (continued):

```
public class RecursiveList<E> implements List<E> {  
    ...  
    public int size() {  
        return size(this.head);  
    }  
    public static <E> int size(ListNode<E> node) {  
        if (node == null) { return 0; }  
        return 1+size(node.getNext());  
    }  
    ...  
}
```

# List ADT: Implementation 5

- Implementation 5 (continued):

```
public class RecursiveList<E> implements List<E> {  
    ...  
    public void add(E elem) {  
        this.head = add(this.head, elem);  
    }  
    public static <E> ListNode<E> add(ListNode<E> n, E e) {  
        if (n == null) { return new ListNode<E>(e, null); }  
        n.setNext(add(n.getNext(), e));  
        return n;  
    }  
    ...  
}
```

# List ADT: Implementation 5

- Implementation 5 (continued):

```
public class RecursiveList<E> implements List<E> { ...
    public void add(int i, E elem) {
        if (i < 0) { throw new IllegalArgumentException(); }
        this.head = add(this.head, i, elem);
    }
    public static <E> ListNode<E> add(ListNode<E> n, int i, E e) {
        if (i == 0) { return new ListNode<E>(e, n); }
        if (n == null) { throw new Index...Exception(); }
        n.setNext(add(n.getNext(), i-1, e));
        return n;
    } ...
}
```

# List ADT: Implementation 5

- Implementation 5 (continued):

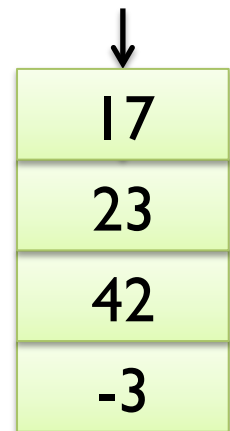
```
public class RecursiveList<E> implements List<E> { ...
    public void remove(int i) {
        if (i < 0) { throw new IllegalArgumentException(); }
        this.head = remove(this.head, i);
    }
    public static <E> ListNode<E> remove(ListNode<E> n, int i) {
        if (n == null) { throw new Index...Exception(); }
        if (i == 0) { return n.getNext(); }
        n.setNext(remove(n.getNext(), i-1));
        return n;
    } } // DONE
```



# **ABSTRACT DATA TYPES FOR STACKS & QUEUES**

# Stacks

- stacks are special sequences, where elements are only added and removed at one end
- imagine a stack of paper on a desk
- many uses:
  - postfix calculator
  - activation records
  - depth-first tree traversals
  - ...
- basic stack operations are
  - looking at the top of the stack
  - removing the top-most element
  - adding an element to the top of the stack



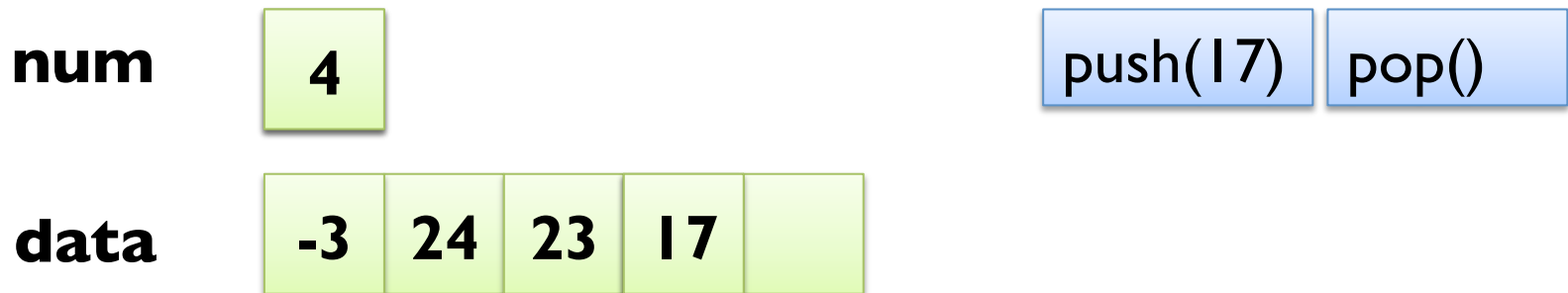
# Stack ADT: Specification

- data are arbitrary objects of class E
- operations are defined by the following interface

```
public interface Stack<E> {  
    public boolean isEmpty();           // is stack empty?  
    public E peek();                   // look at top element  
    public E pop();                     // remove top element  
    public void push(E elem);          // add top element  
}
```

# Stack ADT: Design I

- Design I: use dynamic array
  - the top of the stack is the end of the list
  - in other words, num specifies the top position
  - pushing corresponds to adding at the end
  - popping corresponds to removing at the end



# Stack ADT: Implementation I

- Implementation I:

```
public class DynamicArrayStack<E> implements Stack<E> {  
    private int limit;           // maximal number of elements  
    private E[] data;           // elements of the list  
    private int num = 0;        // current number of elements  
    public DynamicArrayStack(int limit) {  
        this.limit = limit;  
        this.data = (E[]) new Object[limit];  
    }  
    public boolean isEmpty() { return this.num == 0; }  
    ...  
}
```

# Stack ADT: Implementation I

- Implementation I (continued):

```
public class DynamicArrayStack<E> implements Stack<E> { ...
    public E peek() {
        if (this.isEmpty()) { throw new RuntimeException("es"); }
        return this.data[this.num-1];
    }
    public E pop() {
        E result = this.peek();
        num--;
        return result;
    } ...
}
```

# Stack ADT: Implementation I

- Implementation I (continued):

```
public class DynamicArrayStack<E> implements Stack<E> { ...
    public void push(E elem) {
        if (this.num >= this.limit) {
            E[] newData = (E[]) new Object[2*this.limit];
            for (int j = 0; j < limit; j++) { newData[j] = data[j]; }
            this.data = newData;
            this.limit *= 2;
        }
        this.data[num++] = elem;
    }
}
```

# Stack ADT: Design & Implement. 2

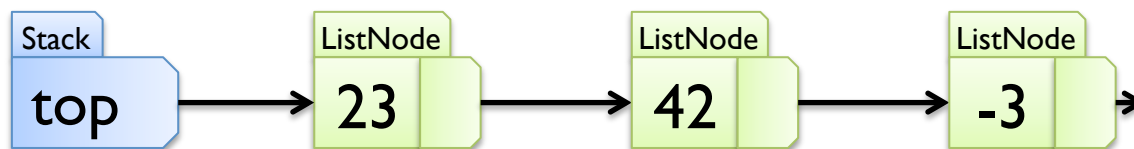
- Design 2: reuse dynamic array list (`ArrayList<E>`)
- Implementation 2:

```
public class ArrayListStack<E> implements Stack<E> {  
    private List<E> list = new ArrayList<E>();  
    public boolean isEmpty() { return this.list.isEmpty(); }  
    public E peek() { return this.list.get(this.list.size()-1); }  
    public E pop() { return this.list.remove(this.list.size()-1); }  
    public void push(E elem) { this.list.add(elem); }  
}
```



# Stack ADT: Design 3

- Design 3: use recursive data structure
  - linked lists have cheap insert and remove operations
  - adding at the end requires running to the end
  - represent top as the beginning of the “list”
- reuse linked list node class (`ListNode<E>`)
- with dynamic arrays, sometimes need to copy full array
- with linked list, always constant time operations



# Stack ADT: Implementation 3

- Implementation 3:

```
public class LinkedStack<E> implements Stack<E> {  
    private ListNode<E> top = null; // top of the stack  
    public boolean isEmpty() { return this.top == null; }  
    public E peek() {  
        if (this.isEmpty()) { throw new RuntimeException("es"); }  
        return this.top.get(0);  
    }  
    ...  
}
```

# Stack ADT: Implementation 3

- Implementation 3 (continued):

```
public class LinkedStack<E> implements Stack<E> {
```

```
...
```

```
public E pop() {
```

```
    E result = this.peek();
```

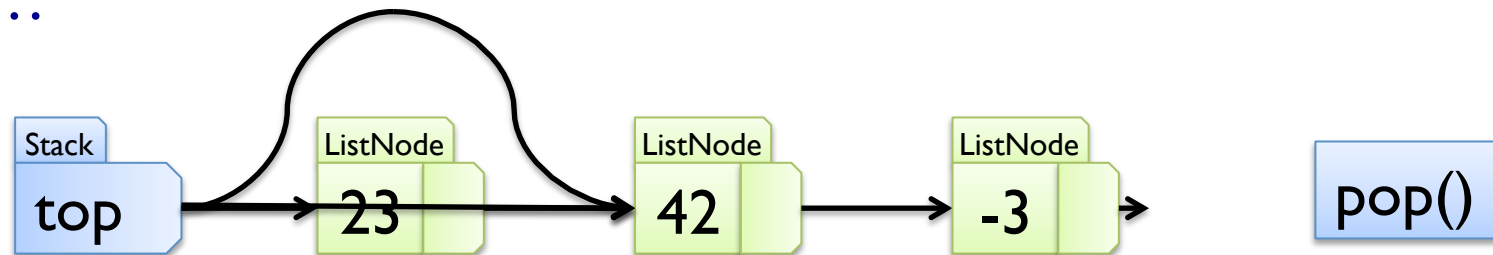
```
    this.top = this.top.getNext();
```

```
    return result;
```

```
}
```

```
...
```

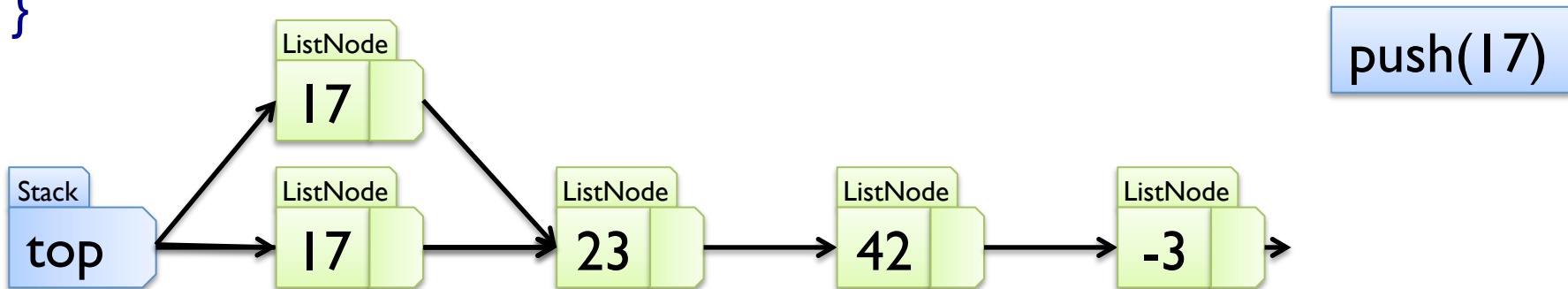
```
}
```



# Stack ADT: Implementation 3

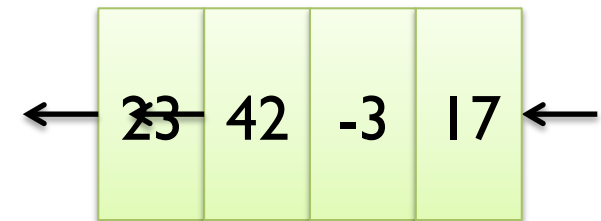
- Implementation 3 (continued):

```
public class LinkedStack<E> implements Stack<E> {  
    private ListNode<E> top = null; // top of the stack  
    ...  
    public void push(E elem) {  
        this.top = new ListNode<E>(elem, this.top);  
    }  
}
```



# Queues

- queues are special sequences, where elements are added on one and removed at the other end
- imagine a waiting line in the supermarket
- many uses:
  - network send/receive buffers
  - process scheduling
  - breadth-first tree traversals
  - ...
- basic queue operations are
  - looking at the beginning of the queue
  - removing the first element
  - adding an element to the end of the queue



# Queue ADT: Specification

- data are arbitrary objects of class E
- operations are defined by the following interface

```
public interface Queue<E> {  
    public boolean isEmpty();           // is queue empty?  
    public E peek();                   // look at first element  
    public E poll();                   // remove first element  
    public boolean offer(E elem);      // true, if element added  
                                        // at end of queue; false, if queue is full  
}
```

# Queue ADT: Design & Implement. I

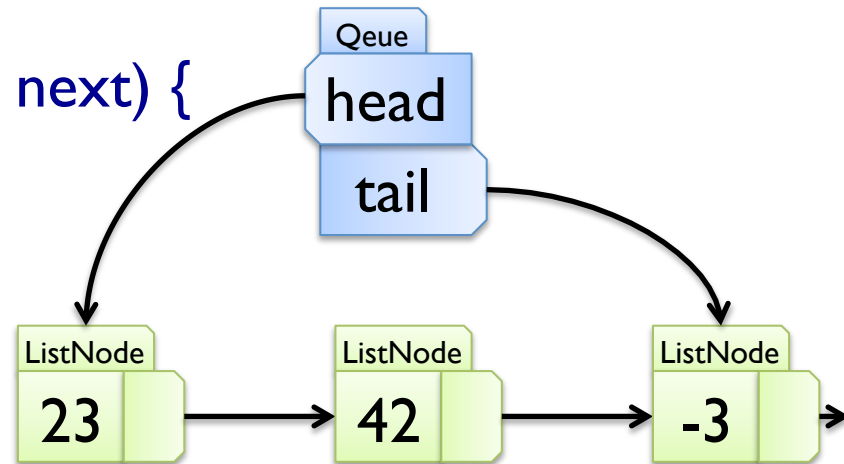
- Design I: reuse dynamic array list (`ArrayList<E>`)
- Implementation I:

```
public class ArrayListQueue<E> implements Queue<E> {  
    private List<E> list = new ArrayList<E>();  
    public boolean isEmpty() { return this.list.isEmpty(); }  
    public E peek() { return this.list.get(0); }  
    public E poll() { return this.list.remove(0); }  
    public boolean offer(E elem) {  
        this.list.add(elem);  
        return true;  
    }  
}
```

# Queue ADT: Design & Implement. 2

- Design 2: use recursive data structure
  - use two references instead of one
  - one reference to end of queue
  - one reference to beginning of queue
- reuse & extend linked list node class (`ListNode<E>`)
- Implementation 2:

```
public class ListNode<E> { ...  
    public void setNext(ListNode<E> next) {  
        this.next = next;  
    }  
}
```





# Queue ADT: Implementation 2

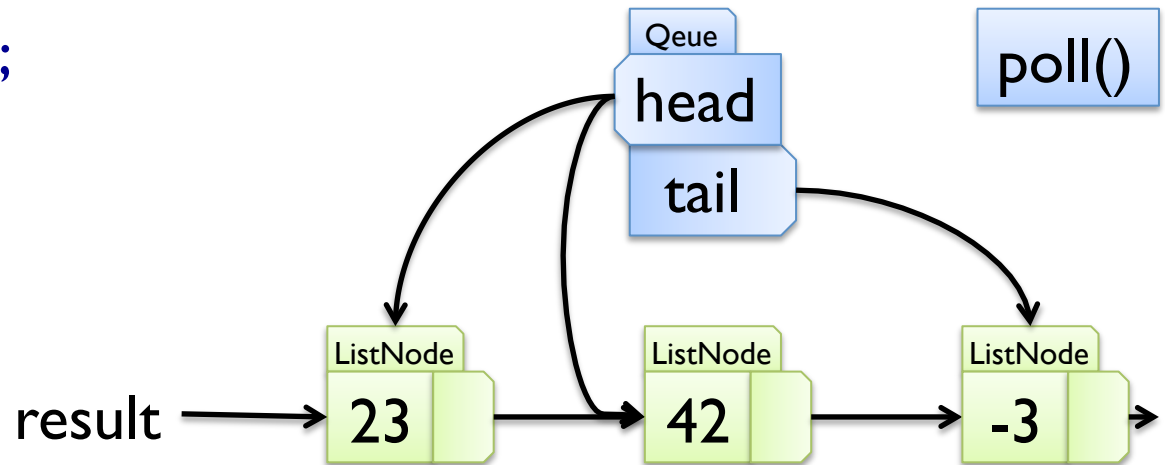
- Implementation 2 (continued):

```
public class LinkedListQueue<E> implements Queue<E> {  
    private ListNode<E> head = null;           // beginning  
    private ListNode<E> tail = null;          // end  
    public boolean isEmpty() {  
        return this.head == null;  
    }  
    public E peek() {  
        return this.head.get(0);  
    }  
    ...  
}
```

# Queue ADT: Implementation 2

- Implementation 2 (continued):

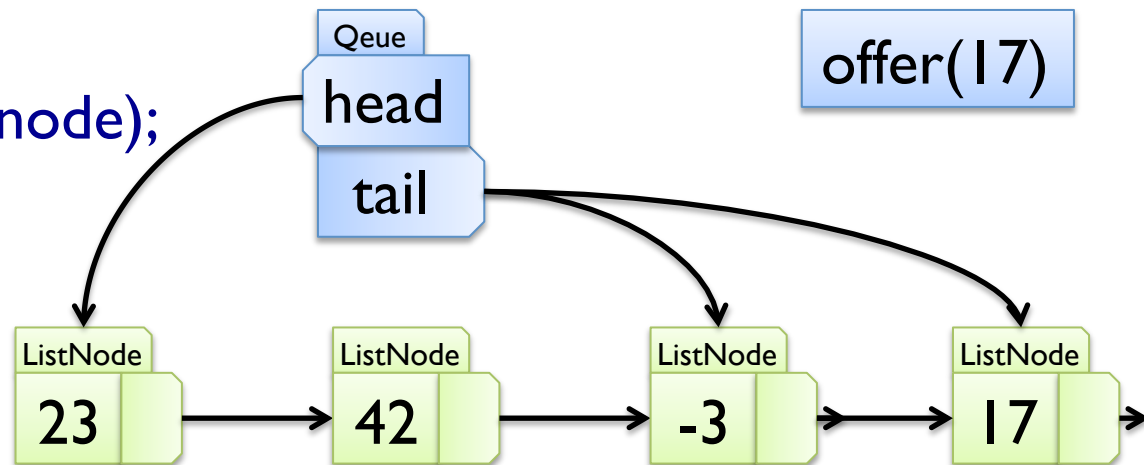
```
public class LinkedQueue<E> implements Queue<E> { ...  
    public E poll() {  
        E result = this.peek();  
        this.head = this.head.getNext();  
        if (this.head == null) {  
            this.tail = null;  
        }  
        return result;  
    }  
    ...  
}
```



# Queue ADT: Implementation 2

- Implementation 2 (continued):

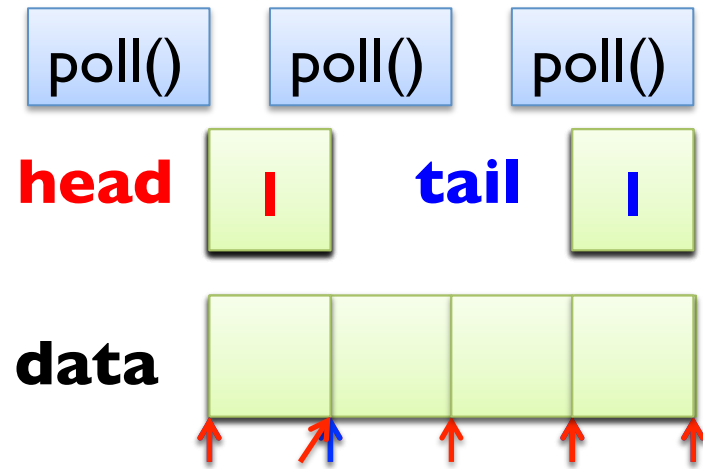
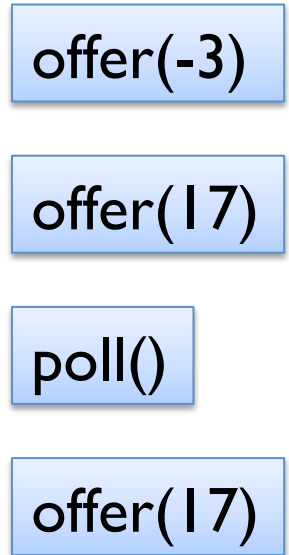
```
public class LinkedQueue<E> implements Queue<E> { ...  
    public boolean offer(E elem) {  
        ListNode<E> node = new ListNode<E>(elem, null);  
        if (this.head == null) {  
            this.head = this.tail = node;  
        } else {  
            this.tail.setNext(node);  
            this.tail = node;  
        }  
        return true;  
    }  
}
```



# Queue ADT: Design & Implement. 3

- Design 3: use a fixed length array
  - use two indices denoting beginning and end
  - wrap around end of array
- very efficient (memory and runtime – no objects!)
- Implementation 3:

```
public class RingQueue<E> implements Queue<E> {
    private int limit;
    private int head = 0; // beginning
    private int tail = 0; // end
    private E[] data;
    ...
}
```



# Queue ADT: Implementation 3

- Implementation 3 (continued):

```
public class RingQueue<E> implements Queue<E> { ...
    private int head = 0;    // beginning
    private int tail = 0;    // end
    private E[] data;
    public boolean isEmpty() { return this.head == this.tail; }
    public E peek() {
        if (this.isEmpty()) { throw new RuntimeException("eq"); }
        return this.data[this.head];
    }
    ...
}
```

# Queue ADT: Implementation 3

- Implementation 3 (continued):

```
public class RingQueue<E> implements Queue<E> {  
    ...  
    public E poll() {  
        E result = this.peek();  
        this.head = (this.head+1) % this.limit;  
        return result;  
    }  
    ...  
}
```

# Queue ADT: Implementation 3

- Implementation 3 (continued):

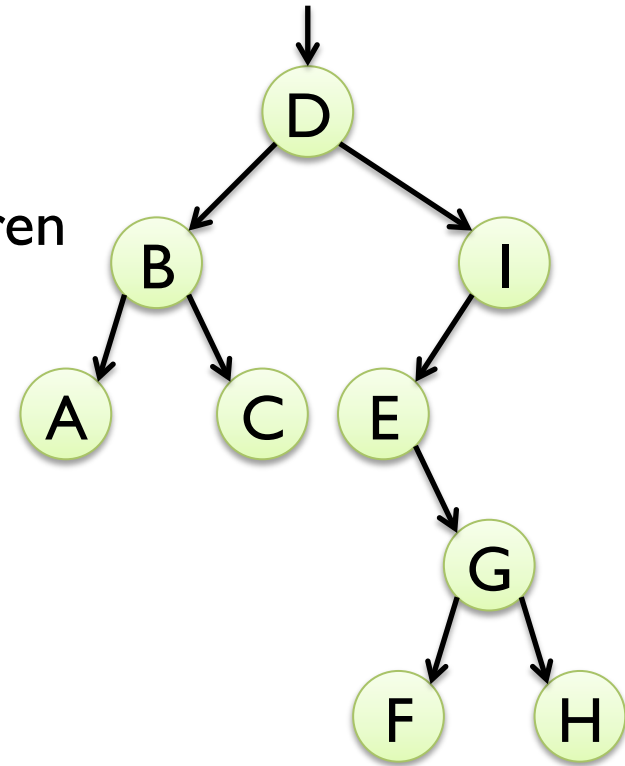
```
public class RingQueue<E> implements Queue<E> { ...
    public boolean offer(E elem) {
        int newTail = (this.tail+1) % this.limit;
        if (newTail == this.head) {
            return false;           // full
        }
        this.data[this.tail] = elem;
        this.tail = newTail;
        return true;
    } ...
}
```

# **ABSTRACT DATA TYPES FOR (BINARY) TREES**



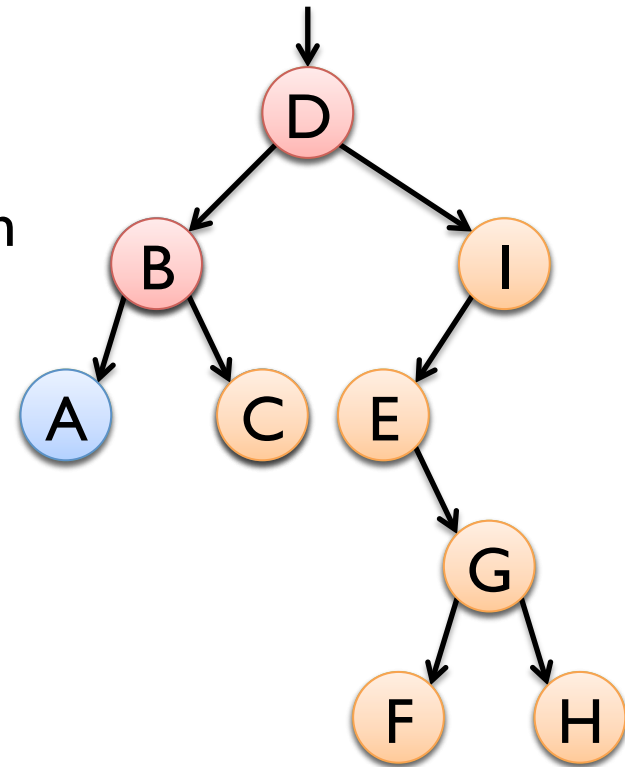
# Trees

- trees store elements non-sequentially
- every node in a tree has 0 or more children
- imagine a tree with root in the air 😊
- many uses:
  - decision tree
  - binary sort trees
  - data base indices
  - ...
- no consensus on what basic binary tree operations are 😞
- set of operations depends on application
- **here:** keeping elements sorted, indexing a database



# Binary Trees

- special case of general trees
- every node in a tree has 0, 1 or 2 children
- tree on the right is an example
- notation:
  - first node is called “**root**”
  - other nodes either in “**left subtree**”
  - ... or in “**right subtree**”
- every node is root in its own subtree!
- for example, look at node B
- node A is the “left child” of B
- node C is the “right child” of B
- node B is the “parent” of both A and C



# BinTree ADT: Specification I

- data are arbitrary objects of class E
- operations are defined by the following interface

```
public interface BinTree<E> {  
    public boolean isEmpty();           // is tree empty?  
    public int size();                 // number of elements  
    public int height();              // maximal depth  
    public boolean contains(E elem);  // true, if elem in tree  
    public List<E> preOrder();        // pre-order traversal  
    public List<E> inOrder();         // in-order traversal  
    public List<E> postOrder();       // post-order traversal  
}
```

# BinTree ADT: Design & Implement. I

- Design I: use recursive data structure
  - based on representing tree nodes by `BinTreeNode<E>`
- Implementation I:

```
public class BinTreeNode<E> {  
    public E elem;  
    public BinTreeNode<E> left, right;  
    public BinTreeNode(E elem, BinTreeNode<E> left,  
                        BinTreeNode<E> right) {  
        this.elem = elem;  
        this.left = left; this.right = right;  
    }  
}
```

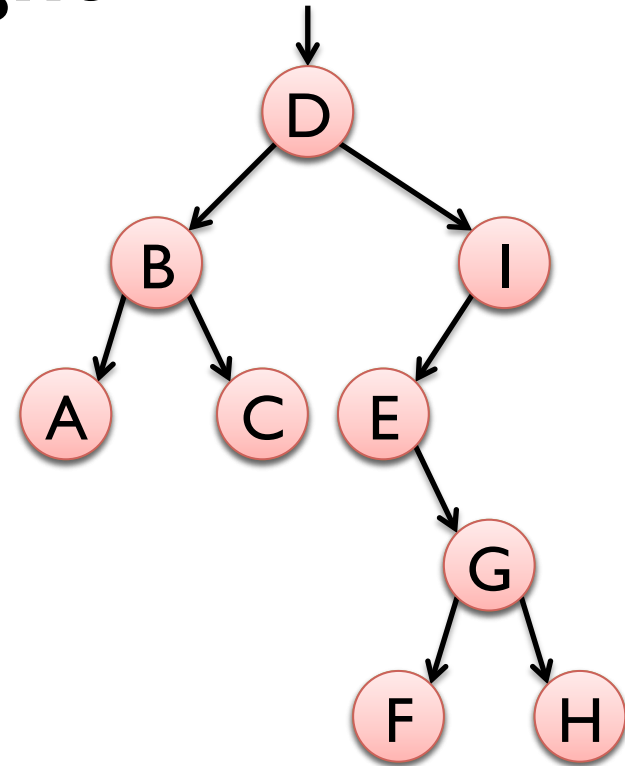
# BinTree ADT: Implementation I

- Implementation I (continued):

```
public class RecursiveBinTree<E> {  
    private BinTreeNode<E> root = null;  
    public int size() { return size(this.root); }  
    private static int size(BinTreeNode<E> node) {  
        if (node == null) { return 0; }  
        return 1 + size(node.left) + size(node.right);  
    }  
    ...  
}
```

# Depth and Height

- depth of the root is 0
- depth of other nodes is  $1 + \text{depth}(\text{parent})$
- Example:
  - 0
  - 1
  - 2
  - 3
  - 4
- height of a subtree is maximal depth of any of its nodes
- Example: height of tree (=subtree starting in D) is 4



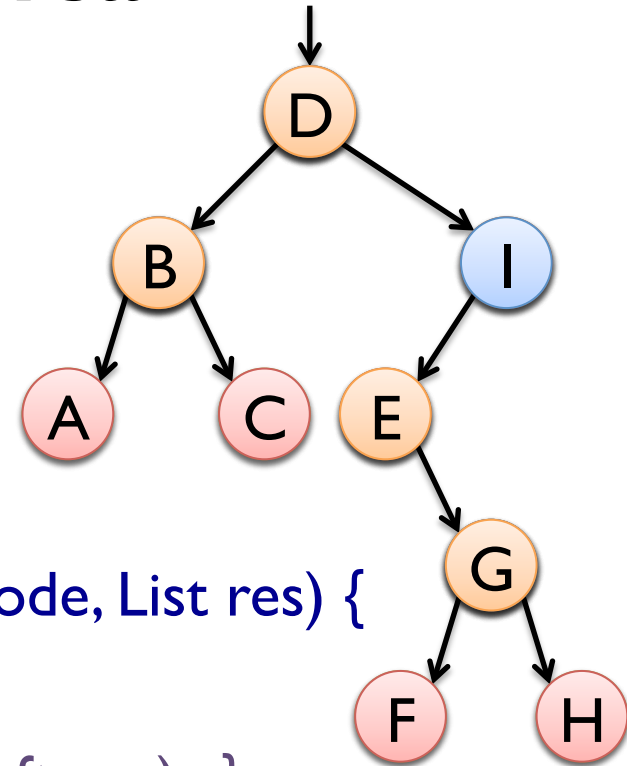
# BinTree ADT: Implementation I

- Implementation I (continued):

```
public class RecursiveBinTree<E> {  
    private BinTreeNode<E> root = null;  
    ...  
    public int height() { return height(this.root); }  
    private static int height(BinTreeNode<E> node) {  
        if (node == null) { return -1; }  
        return 1 + max(height(node.left), height(node.right));  
    }  
    private static int max(int a, int b) { return a > b ? a : b; }  
    ...  
}
```

# Binary Tree Traversal

- traversal can be either
  - depth-first
  - breadth-first
- three standard depth-first traversals
  - pre-order



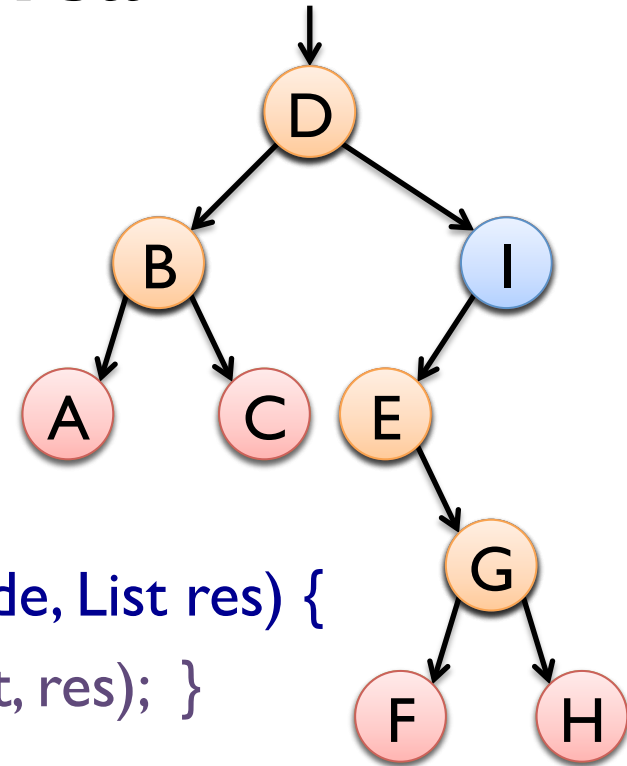
```
public static void preOrder(BinTreeNode node, List res) {  
    res.add(node.elem);  
    if (node.left != null) { preOrder(node.left, res); }  
    if (node.right != null) { preOrder(node.right, res); }  
}
```





# Binary Tree Traversal

- traversal can be either
  - depth-first
  - breadth-first
- three standard depth-first traversals
  1. pre-order
  2. in-order
  3. post-order

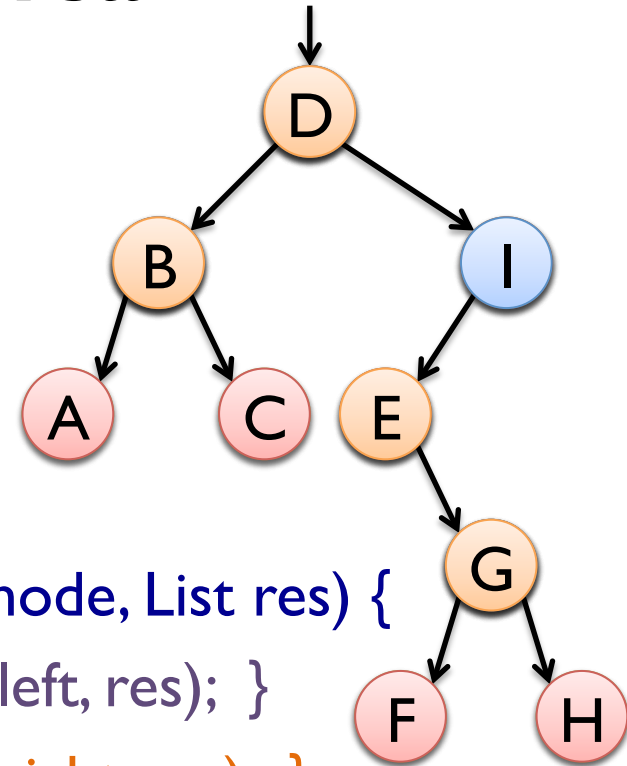


```
public static void inOrder(BinTreeNode node, List res) {  
    if (node.left != null) { inOrder(node.left, res); }  
    res.add(node.elem);  
    if (node.right != null) { inOrder(node.right, res); }  
}
```

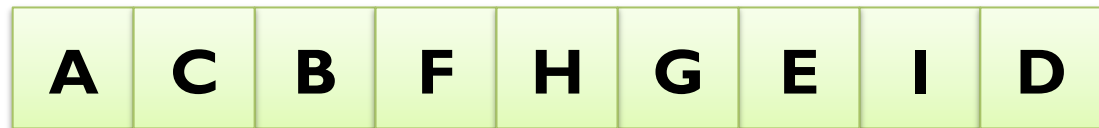


# Binary Tree Traversal

- traversal can be either
  - depth-first
  - breadth-first
- three standard depth-first traversals
  - pre-order
  - in-order
  3. post-order



```
public static void postOrder(BinTreeNode node, List res) {  
    if (node.left != null) { postOrder(node.left, res); }  
    if (node.right != null) { postOrder(node.right, res); }  
    res.add(node.elem);  
}
```



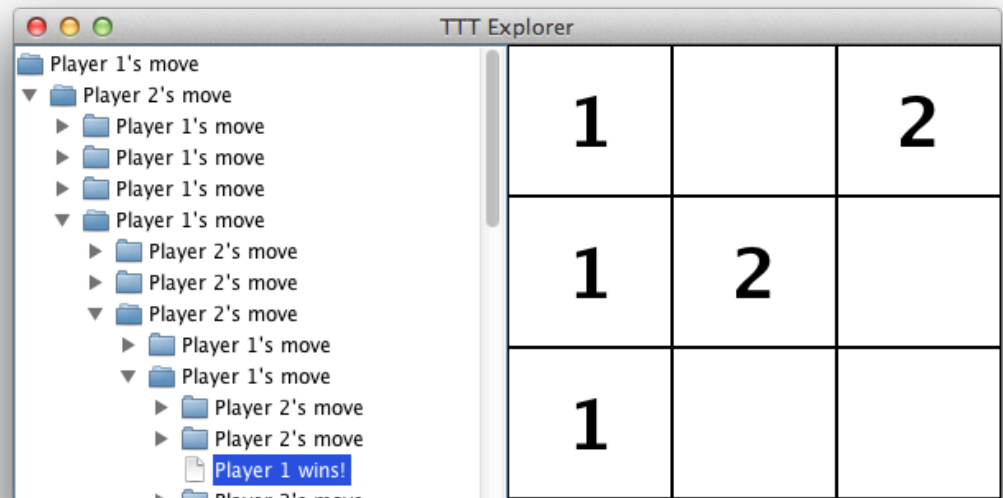
# PROJECT PART 2

# Organizational Details

- exam project consisting of 2 parts
- both parts have to be passed to pass the course
- projects must be done individually, so no co-operation
- you may talk about the problem and ideas how to solve them
- deliverables:
  - written 4 page report as specified in project description
  - handed in as a single PDF file
  - deadline: Friday, January 9, 2013, 12:00
- ENOUGH - now for the ABSTRACT part ...

# Board Games: Tic Tac Toe & Co

- n-way Tic Tac Toe:
  - n player board game played on  $(n+1) \times (n+1)$  grid
  - first player to place 3 marks in a row/column/diagonal wins
- **Goal:** implement ADT for game tree (viewer)
- **Challenges:**
  - ADT Design
  - ADT Implementation
  - Multivariate Trees



# Board Games: Tic Tac Toe & Co

- Task 0: Preparation
  - download and understand existing framework
  - integrate (and fix) **TTTBoard** and **Coordinate**
- Task 1: Implement ADT
  - design and implement **TTTGameTree** (and node class)
  - need to cooperate with **GameTreeDisplay** and **TTTExplorer**
- Task 2: Building the Game Tree
  - build a tree with one node representing the initial game
  - add successors of a game state as children
  - keep going until one player wins or a draw is reached
  - check you progress using **TTTExplorer**

# Board Games: Tic Tac Toe & Co

- Task 3 (optional): Reducing the Size of the Game Tree
  - reuse nodes for identical game states
  - consider rotational symmetry?
  - consider mirroring?
- Task 4 (optional): Artificial Intelligence
  - fully expanded game tree useful to implement AI player
  - AI player tries to develop towards winning situations
  - AI player tries to avoid losing situations