

Rosen 8.1 Recurrence relations

A recurrence relation $\{a_n\}$ is an equation that expresses a_n in terms of $a_0, a_1, \dots, a_{n-1}$ for all $n \geq n_0$, where $n_0 \geq 0$

$\{x_n\}$ is a solution to $\{a_n\}$ if each term x_i satisfies the same equation as a_i

Example: $a_n = 3n$ is a solution to

$$a_n = 2a_{n-1} - a_{n-2} \text{ for all } n = 2, 3, \dots$$

$$\text{as } 2 \cdot 3(n-1) - 3(n-2)$$

$$= 6n - 6 - 3n + 6 = 3n$$

Initial conditions: Then specify the first terms when the recurrence starts (up to n_0)

Above we could take $a_0 = 0$ and $a_1 = 3$ and see that

$$\text{then } 6 = 3 \cdot 2 = a_2 = 2a_1 - a_0 = 2 \cdot 3 - 0 = 6$$

$$\text{and } 9 = 3 \cdot 3 = a_3 = 2a_2 - a_1 = 2 \cdot 6 - 3 = 9$$

Initial conditions uniquely determine the sequence (induction on n)

Example 1 Fibonacci numbers

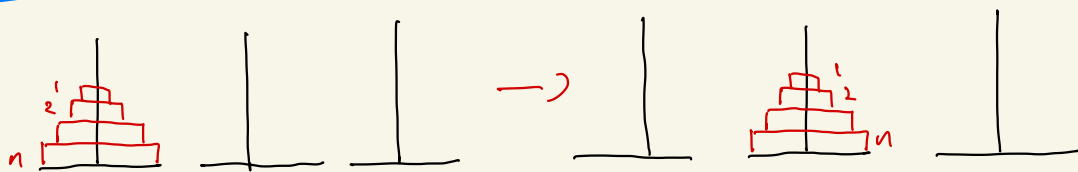
$$f_1 = f_2 = 1$$

$$f_n = f_{n-1} + f_{n-2}$$

$$f_3 = f_2 + f_1 = 1 + 1 = 2, \quad f_4 = f_3 + f_2 = 2 + 1 = 3$$

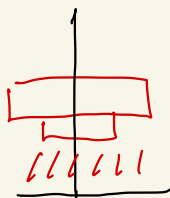
$$f_5 = 5, \quad f_6 = 8 \quad \dots$$

Example 2 Tower of Hanoi Puzzle



Initial

rule never



$H_n = \# \text{ moves when there are } n \text{ discs}$

$$H_1 = 1$$

$$H_n \leq 2H_{n-1} + 1$$

1. $n-1$ top discs $\rightarrow$ peg no 3

2. move disc $n \rightarrow$ peg 2

3. move $n-1$ discs peg 3 $\rightarrow$ peg 2

$$H_n \geq 2H_{n-1} + 1$$

- Cannot move disc n before all discs $1, 2, \dots, n-1$ have been moved to another peg and this must be peg 3 as we want disc n on peg 2

Thus we need at least H_{n-1} moves for this part

- 1 move to get disc n to peg 2

- H_{n-1} moves to get discs $1, 2, \dots, n-1$ from peg 3 to peg 2. Corresponds to problem with n replaced by $n-1$

Conclusion: $H_n = 2H_{n-1} + 1$ (and $H_1 = 1$)

$$\begin{aligned} H_n = 2^n - 1 : \quad 2^n - 1 &= H_n = 2H_{n-1} + 1 \\ &= 2(2^{n-1} - 1) + 1 \\ &= 2^n - 1 \end{aligned}$$

We will see a method to solve such recurrence equations in Section 8.2

Example 3 # of bitstrings of length n with no 00

Let N be the # of such strings

By some rule, $N = N_1 + N_0$ where N_i is # of such strings that end with i $i=0,1$

Then we see that

$$N_1 = \# \text{ length } n-1 \text{ with no } \sim 00 \sim$$

$$N_0 = \# \text{ length } n-2 \text{ with no } \sim 00 \sim$$

(since bit $n-1$ must be a 1 ~ 10)

$$\text{Hence } a_n = N = N_1 + N_0 = a_{n-1} + a_{n-2} \text{ when } n \geq 3$$

$$a_n = a_{n-1} + a_{n-2}$$

$$\text{initial conditions } a_1 = 2 \text{ ('0', '1')} \quad \text{so } a_1 = f_3$$

$$a_2 = 3 \text{ ('01', '10', '11')} \quad \text{so } a_2 = f_4$$

$$a_3 = a_2 + a_1 = f_4 + f_3 = f_5$$

$$a_4 = a_3 + a_2 = f_5 + f_4 = f_6$$

$$\text{So } a_n = f_{n+2} \text{ for } n \geq 1$$

Thm 8.2 Solving linear recurrence equations

Def 1 A linear homogeneous recurrence relation of degree k with constant coefficients is of the form

$$(*) \quad a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} \quad \text{when } c_i \in \mathbb{R} \text{ and } a_k \neq 0$$

Note the solution to $(*)$ is uniquely determined by the values of $a_0, a_1, \dots, a_{k-1}$ (induction)

$(*)$ is homogeneous because a_n only depends on previous terms a_i and not on some function $f(n)$

$(*)$ has degree k as a_n is expressed as a linear combination of terms whose index is at least k .

How to solve $(*)$?

Make a guess $a_n = r^n$ for some $r \neq 0$ and insert

we must have

$$r^n = c_1 r^{n-1} + c_2 r^{n-2} + \dots + c_k r^{n-k}$$

So

$$r^k = c_1 r^{k-1} + c_2 r^{k-2} + \dots + c_k$$

or

$$r^k - c_1 r^{k-1} - c_2 r^{k-2} - \dots - c_k = 0$$

divided by r^{n-k}

$$(D) \quad r^k - C_1 r^{k-1} - C_2 r^{k-2} - \dots - C_k = 0$$

is the characteristic equation for the sequence $\{a_n\}$

satisfying $a_n = C_1 a_{n-1} + C_2 a_{n-2} + \dots + C_k a_{n-k}$ (*)

Solutions to (D) are called characteristic roots.

Lemma Suppose $\{s_n\}$ and $\{t_n\}$ are solutions to
and c is a constant then

1) $\{c \cdot s_n\}$ is a solution to (*)

2) $\forall c, d$ constants $\{c s_n + d t_n\}$ is a solution to (*)

proof:

$$s_n = C_1 s_{n-1} + C_2 s_{n-2} + \dots + C_k s_{n-k}$$

$\Downarrow$

$$c \cdot s_n = c (C_1 s_{n-1} + C_2 s_{n-2} + \dots + C_k s_{n-k})$$

$$= C_1 \underbrace{c s_{n-1}} + C_2 \underbrace{c s_{n-2}} + \dots + C_k \underbrace{c s_{n-k}}$$

so $\{c s_n\}$ solves *

if $s_n = C_1 s_{n-1} + C_2 s_{n-2} + \dots + C_k s_{n-k}$

and $t_n = C_1 t_{n-1} + C_2 t_{n-2} + \dots + C_k t_{n-k}$

then $s_n + t_n = C_1 (s_{n-1} + t_{n-1}) + C_2 (s_{n-2} + t_{n-2}) + \dots + C_k (s_{n-k} + t_{n-k})$

Now 2) follows from this and 1)

□.

The case when $k=2$

$$(*) \quad a_n = c_1 a_{n-1} + c_2 a_{n-2} \sim r^2 - c_1 r - c_2 = 0 \text{ (characteristic eq.)}$$

Theorem 1 Suppose $r^2 - c_1 r - c_2 = 0$ has distinct roots r_1, r_2

Then $\{a_n\}$ solves $*$ if and only if

$$a_n = \alpha_1 r_1^n + \alpha_2 r_2^n \quad \text{for } n=1, 2, \dots \text{ and } \alpha_1, \alpha_2 \in \mathbb{R}$$

Proof: let $r_1 \neq r_2$ be roots in $r^2 - c_1 r - c_2 = 0$

then $r_i^2 - c_1 r_i - c_2 = 0 \Rightarrow r_i^2 = c_1 r_i + c_2$ for $i=1, 2$ (A)

Now if we set $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n$ for some $\alpha_1, \alpha_2 \in \mathbb{R}$

we get

$$\begin{aligned} c_1 a_{n-1} + c_2 a_{n-2} &= c_1 (\alpha_1 r_1^{n-1} + \alpha_2 r_2^{n-1}) + c_2 (\alpha_1 r_1^{n-2} + \alpha_2 r_2^{n-2}) \\ &= [c_1 \alpha_1 r_1^{n-1} + c_2 \alpha_1 r_1^{n-2}] + [c_1 \alpha_2 r_2^{n-1} + c_2 \alpha_2 r_2^{n-2}] \\ &= \alpha_1 r_1^{n-2} [c_1 r_1 + c_2] + \alpha_2 r_2^{n-2} [c_1 r_2 + c_2] \quad \text{by (A)} \\ &= \alpha_1 r_1^{n-2} \cdot r_1^2 + \alpha_2 r_2^{n-2} \cdot r_2^2 \\ &= \alpha_1 r_1^n + \alpha_2 r_2^n \end{aligned}$$

Hence $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n$ is a solution to $(*)$

Conversely, suppose $\{a_n\}$ solves (*) with

$$a_0 = C_0 \text{ and } a_1 = C_1$$

This uniquely determines $\{a_n\}$ so we need to find α_1, α_2 such that $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n$

$$\text{with } a_0 = C_0, a_1 = C_1$$

we get 2 equations with 2 unknown α_1, α_2 :

$$C_0 = a_0 = \alpha_1 \cdot r_1^0 + \alpha_2 \cdot r_2^0 = \alpha_1 + \alpha_2$$

$$C_1 = a_1 = \alpha_1 r_1 + \alpha_2 r_2$$

$$C_0 = \alpha_1 + \alpha_2$$

$$C_1 = \alpha_1 r_1 + \alpha_2 r_2$$

$\Rightarrow$

$$C_1 - C_0 r_2 = \alpha_1 r_1 - \alpha_1 r_2$$

$$\text{so } \boxed{\alpha_1 = \frac{C_1 - C_0 r_2}{r_1 - r_2}}$$

$\Rightarrow$

$$C_1 - C_0 r_1 = \alpha_2 r_2 - \alpha_2 r_1$$

$$\text{so } \boxed{\alpha_2 = \frac{C_1 - C_0 r_1}{r_2 - r_1}}$$

Hence there is a unique choice of α_1, α_2 such that $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n$ solves

(*) with initial conditions $a_0 = C_0$
 $a_1 = C_1$

Example 3

$$a_n = a_{n-1} + 2a_{n-2}$$

$$a_0 = 2, a_1 = 7$$

Characteristic equation: $r^2 - r - 2 = 0$

roots $r_1 = 2, r_2 = -1$

$$\text{so } a_n = \alpha_1 2^n + \alpha_2 (-1)^n$$

Determining α_1, α_2 :

$$\alpha_1 = \frac{C_1 - C_0 r_2}{r_1 - r_2} = \frac{7 - 2 \cdot (-1)}{2 - (-1)} = \frac{9}{3} = 3$$

$$\alpha_2 = \frac{C_1 - C_0 r_1}{r_2 - r_1} = \frac{7 - 2 \cdot 2}{-1 - 2} = -1$$

$$a_n = 3 \cdot 2^n - (-1)^n$$

Theorem 2 Let $C_1, C_2 \in \mathbb{R}$ with $C_2 \neq 0$ and
suppose $r^2 - C_1 r - C_2 = 0$ has a unique (double) root r
Then $\{a_n\}$ solves $a_n = C_1 a_{n-1} + C_2 a_{n-2}$ if and only if

$$a_n = \alpha_1 \cdot r^n + \alpha_2 \cdot n \cdot r^n \text{ for some } \alpha_1, \alpha_2 \in \mathbb{R}$$

proof: exercis!

Example 5 $a_n = 6a_{n-1} - 9a_{n-2}$ $a_0 = 1$ $a_1 = 6$

Characteristic eq: $r^2 - 6r + 9 = 0 \Rightarrow r = 3$

So $a_n = \alpha_1 3^n + \alpha_2 n \cdot 3^n$

Find α_1, α_2 : $C_0 = 1 = a_0 = \alpha_1 \cdot 3^0 + \alpha_2 \cdot 0 \cdot 3^0 = \alpha_1$

$$\begin{aligned} C_1 = 6 = a_1 &= \alpha_1 \cdot 3^1 + \alpha_2 \cdot 1 \cdot 3^1 \\ &= 3\alpha_1 + 3\alpha_2 \\ &= 3 + 3\alpha_2 \text{ so } \alpha_2 = 1 \end{aligned}$$

so $\alpha_2 = (6-3)/3 = 1$

$$a_n = 3^n + n \cdot 3^n = (n+1)3^n$$

Theorem 3 $c_i \in \mathbb{R}$ $i=1, 2, \dots, k$, $c_k \neq 0$

if $r^k - c_1 r^{k-1} - c_2 r^{k-2} - \dots - c_k = 0$ has distinct roots

$r_1, r_2, \dots, r_k$ then $\{a_n\}$ solves

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$$

if and only if

$$a_n = \alpha_1 r_1^n + \alpha_2 r_2^n + \dots + \alpha_k r_k^n$$

proof exercis

8.2.3 Linear inhomogeneous recurrence relations with constant coefficients

Example of such a recurrence relation $a_n = 3a_{n-1} + 2n$

generally $a_n = C_1 a_{n-1} + C_2 a_{n-2} + \dots + C_k a_{n-k} + F(n)$ (*)

inhomogeneous because of the term $F(n)$

Associated homogeneous equation to (*) is

$$a_n = C_1 a_{n-1} + C_2 a_{n-2} + \dots + C_k a_{n-k} \quad (**)$$

Key fact: Every solution to (*) can be obtained from any (fixed) solution

Theorem 5 Let $\{a_n^P\}$ be any solution to (*) (called a particular solution). Then every solution to (*) is of the form $\{a_n^h + a_n^P\}$ where $\{a_n^h\}$ is a solution to (**)

Proof Suppose $\{b_n\}$ also solves (*) (besides $\{a_n^P\}$). Then

$$b_n - a_n^P = C_1(b_{n-1} - a_{n-1}^P) + C_2(b_{n-2} - a_{n-2}^P) + \dots + C_k(b_{n-k} - a_{n-k}^P) \quad \begin{matrix} F(n) \\ \text{terminals} \end{matrix}$$

So $\{b_n - a_n^P\}$ solves (**) and hence $\{b_n\} = \{a_n^h + a_n^P\}$

□.

So key step when solving $(*)$ is to find one solution and after that solve $(*)$

$$\Sigma \times 10 \quad a_n = 3a_{n-1} + 2n \quad a_1 = 3$$

homogeneous part $a_n = 3a_{n-1}$

Char. eq $r - 3 = 0 \Rightarrow \text{root} = r$

$a_n^h = \alpha 3^n$ homogeneous solution

Finding $\{a_n^p\}$:

a) $F(n) = 2n$ is a linear function we guess a linear function $p_n = cn + d$ as a solution to $(*)$

Then we must have $cn + d = 3(cn-1) + d + 2n$
 $= 3cn - 3c + d + 2n$
 $= (3c+2)n + 3(d-c)$

must hold for all n so $c = 3c+2$ and $d = 3(d-c)$
 $\Downarrow c = -1$ and $d = 3(d+1)$
 $\Downarrow d = -\frac{3}{2}$

so $a_n^p = -n - \frac{3}{2}$ is a solution to $(*)$

By Thm 5, all solutions to $(*)$ are of the form $a_n^h + a_n^p$ and we know $a_n^h = \alpha 3^n$

so $a_n = \alpha \cdot 3^n - n - \frac{3}{2}$ is the general solution

we have $a_1 = 3$ so $3 = \alpha \cdot 3^1 - 1 - \frac{3}{2} = 3\alpha - \frac{5}{2}$
 $\Downarrow 3 = 3\alpha - \frac{5}{2} \Rightarrow \alpha = \frac{11}{6}$

$a_n = \frac{11}{6} \cdot 3^n - n - \frac{3}{2}$

Theorem 6 Suppose $\{a_n\}$ solves

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n) \quad (*)$$

where $F(n) = (b_t n^t + b_{t-1} n^{t-1} + \dots + b_1 n + b_0) S^n$

• If S is NOT a root in the characteristic eq

$$(\square) \quad r^k - c_1 r^{k-1} - c_2 r^{k-2} - \dots - c_k = 0$$

then there is a solution to $(*)$ of the form

$$(p_t n^t + p_{t-1} n^{t-1} + \dots + p_0) \cdot S^n$$

• If S IS a root in $(\square)$ and has multiplicity m

then there is a solution to $(*)$ of the form

$$n^m (p_t n^t + p_{t-1} n^{t-1} + \dots + p_0) \cdot S^n$$

Proof NOT PENSON!

