

## DM42 – Advanced Data Structures

### Addition to Weekly Note 2

#### Defining Asymptotic Notation

Let  $\mathbb{N}$  = denote the natural numbers  $\{0, 1, 2, \dots\}$  and let  $\mathbb{R}^+$  the positive real numbers.

$$O(f) = \{g: \mathbb{N} \rightarrow \mathbb{R}^+ \mid \exists c \in \mathbb{R}^+ \exists n_0 \in \mathbb{N} \forall n \in \mathbb{N}: n \geq n_0 \Rightarrow g(n) \leq cf(n)\}$$

$$\Omega(f) = \{g: \mathbb{N} \rightarrow \mathbb{R}^+ \mid \exists c \in \mathbb{R}^+ \exists n_0 \in \mathbb{N} \forall n \in \mathbb{N}: n \geq n_0 \Rightarrow g(n) \geq cf(n)\}$$

$$\Theta(f) = O(f) \cap \Omega(f)$$

$$o(f) = O(f) \setminus \Theta(f)$$

$$\omega(f) = \Omega(f) \setminus \Theta(f)$$

$$O(f(m, n)) = \{g: \mathbb{N}^2 \rightarrow \mathbb{R}^+ \mid \exists c \in \mathbb{R}^+ \exists m_0, n_0 \in \mathbb{N} \forall m, n \in \mathbb{N}: m \geq m_0 \wedge n \geq n_0 \Rightarrow g(m, n) \leq cf(m, n)\}$$

[there are many alternative ways of defining asymptotic notation]

#### Repetition Problems

1. Show that  $O(\log_a n) = O(\log_b n)$ , where  $a, b > 1$ .
2. Show that  $O(n) \subset O(n \log n) \subset O(n^2)$ .
3. Fill in the following table with X's; and arguments.

$A$	$B$	$A \in O(B)$	$A \in o(B)$	$A \in \Omega(B)$	$A \in \omega(B)$	$A \in \Theta(B)$
$\log \log n$	$\log n$					
$(\log n)^c$	$n^k$					
$\frac{\log n}{\log \log n}$	$\log \log n$					
$\sqrt{n}$	$n^{\sin n}$					
$\log n!$	$\log n^n$					

where  $c$  and  $k$  are positive constants.

4. Let  $c, c_1, c_2$  be constants. How does  $T(n)$  grow asymptotically with the following definitions of  $T$ ?
  - (a)  $T(n) = T(\frac{n}{2}) + c$
  - (b)  $T(n) = 2T(\frac{n}{2}) + c$
  - (c)  $T(n) = 3T(\frac{n}{2}) + c$
  - (d)  $T(n) = T(\frac{n}{2}) + n$
  - (e)  $T(n) = 3T(\frac{n}{2}) + n$
  - (f)  $T(n) = T(n - c_1) + c_2$
  - (g)  $T(n) = T(n - c) + n$

Assume that  $n$  is on some convenient form (a power of two or similar is often helpful) and that  $T(1)$  is some (appropriate) constant.