

DM877  
Discrete Optimization

Lecture 4  
**Introduction to MiniZinc**

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# Outline

1. MiniZinc

# Resume

- ▶ Modelling in MILP and CP
  - ▶ First example: graph labelling with consecutive numbers
  - ▶ Second example: Cryptarithmic (or verbal arithmetic or cryptarithm):  $\text{Send} + \text{More} = \text{Money}$
- ▶ Overview on constraint programming:  
representation (modeling language) + reasoning (search + propagation)
  - ▶ search = backtracking + branching
  - ▶ propagate, filtering, pruning
  - ▶ level of consistency (arc/generalized + value/bound/domain)

# Outline

1. MiniZinc

# MiniZinc

- ▶ language for specifying: constrained optimization and decision problems over integers and real numbers.
- ▶ existential and universal quantifiers, sums over index sets, or logical connectives like implications and if-then-else statements
- ▶ supports defining predicates and functions that let users structure their models
- ▶ models are parametric (instance vs problem)
- ▶ MiniZinc compiler transforms **model file + data file** into FlatZinc model via libraries and predicate definitions that target the specific solvers, such as Constraint Programming (CP), Mixed Integer Linear Programming (MIP) or Boolean Satisfiability (SAT) solvers.  
Eg: MiniZinc allows the specification of **global constraints** by **decomposition**.
- ▶ allows **annotations** of the model to let the user fine tune the behaviour of the solver, independent of the declarative meaning of the model.
- ▶ third-party solvers need an **executable** and a solver specific MiniZinc **library**. They must be specified in a MiniZinc **configuration file**.

# A First Script

```
% Colouring Australia using nc colours
int: nc = 3;
var 1..nc: wa; var 1..nc: nt; var 1..nc: sa; var 1..nc: q;
var 1..nc: nsw; var 1..nc: v; var 1..nc: t;

constraint wa != nt;
constraint wa != sa;
constraint nt != sa;
constraint nt != q;
constraint sa != q;
constraint sa != nsw;
constraint sa != v;
constraint q != nsw;
constraint nsw != v;

solve satisfy;

output ["wa=\(wa)\t nt=\(nt)\t sa=\(sa)\n",
"q=\(q)\t nsw=\(nsw)\t v=\(v)\n",
"t=", show(t), "\n"];
```

# Parameters and Variables

- ▶ comments `%` or `/* */`
- ▶ basic parameter types are integers (`int`), floating point numbers (`float`), Booleans (`bool`) and strings (`string`). Arrays and sets are also supported.
- ▶ `decision variables` are variables in the sense of mathematical or logical variables not in the sense of programming language variables.
- ▶ the `domain` can be given as part of the variable declaration and the type is inferred from there
- ▶ `decision variables` can be Booleans, integers, floating point numbers, or sets. Moreover, arrays of decision variables
- ▶ `identifiers` are made of characters or `_` and must start with a char.

Parameters:

```
int : <var-name>  
<l>..<u> : <var-name>
```

Decision variables:

```
var int : <var-name>  
var <l>..<u> : <var-name>
```

variables instantiation vs type



# Constraints

Relational operators:

equal	= or ==
not equal	!=
strictly less than	<
strictly greater than	>
less than or equal to	<=
greater than or equal to	>=

# Output

- ▶ between quotes for strings and enclosed a show call if MiniZinc elements (or shortcut "(e)").
- ▶ string concatenation via operator ++
- ▶ `show_int(n,X)` and `show_float(n,d,X)`

# Command Line Execution

```
$ minizinc --solver gecode aust.mzn
```

```
$ minizinc -h
```

```
$ minizinc -a --solver gecode aust.mzn
```

# Another Example: Production Planning

```
% Baking cakes for the school fete
var 0..100: b; % no. of banana cakes
var 0..100: c; % no. of chocolate cakes
% flour
constraint 250*b + 200*c <= 4000;
% bananas
constraint 2*b <= 6;
% sugar
constraint 75*b + 150*c <= 2000;
% butter
constraint 100*b + 150*c <= 500;
% cocoa
constraint 75*c <= 500;
% maximize our profit
solve maximize 400*b + 450*c;
output ["no. of banana cakes = \ (b)\n",
       "no. of chocolate cakes = \ (c)\n"];
```

## Arithmetics:

Addition	+	integer division	div
subtraction	-	integer modulus	mod
multiplication	*	absolute value	abs
power function	pow		

# Data Files

```
$ minizinc cakes2.mzn -D "flour=4000;banana=6;sugar=2000;butter=500;cocoa=500;"
```

```
$ minizinc cakes2.mzn pantry.dzn
```

- ▶ solution output is in the same format as the dzn
- ▶ check parameters `assert(B,S)`

```
constraint assert(flour >= 0,"Invalid datafile: " ++ "Amount of flour should be non-negative");
```

# A Real Numbers Example

```
% variables
var float: R; % quarterly repayment
var float: P; % principal initially borrowed
var 0.0 .. 10.0: I; % interest rate

% intermediate variables
var float: B1; % balance after one quarter
var float: B2; % balance after two quarters
var float: B3; % balance after three quarters
var float: B4; % balance owing at end

constraint B1 = P * (1.0 + I) - R;
constraint B2 = B1 * (1.0 + I) - R;
constraint B3 = B2 * (1.0 + I) - R;
constraint B4 = B3 * (1.0 + I) - R;

solve satisfy;

output [
"Borrowing ", show_float(0, 2, P), " at ", show(I*100.0),
"% interest, and repaying ", show_float(0, 2, R),
"\nper quarter for 1 year leaves ", show_float(0, 2, B4), " owing\n"
];
```

- ▶ if I borrow \$1000 at 4% and repay \$260 per quarter, how much do I end up owing?
- ▶ if I want to borrow \$1000 at 4% and owe nothing at the end, how much do I need to repay?
- ▶ if I can repay \$250 a quarter, how much can I borrow at 4% to end up owing nothing?

- ▶ needs appropriate solvers
- ▶ addition (+), subtraction (-), multiplication (\*) and floating point division (/).
- ▶ int2float (implicitly used when  $a / b$  and  $a$  and  $b$  are integers)

▶ arithmetic functions:

absolute value (abs),	sine (sin),	hyperbolic sine (sinh),
square root (sqrt) and power (pow)	cosine (cos),	hyperbolic cosine (cosh),
natural logarithm (ln),	tangent (tan),	hyperbolic tangent (tanh),
logarithm base 2 (log2),	arcsine (asin),	hyperbolic arcsine (asinh),
logarithm base 10 (log10),	arc-cosine (acos),	hyperbolic arccosine (acosh),
exponentiation of e (exp),	arctangent (atan),	hyperbolic arctangent (atanh)

- ▶ . sign for decimal
- ▶ Solve using LP ut only if linear constraints



# Basic Structure

```
include <filename>;  
  
<type inst expr>: <variable> [ = ] <expression>;  
  
<variable> = <expression>;  
  
constraint <Boolean expression>;  
  
solve satisfy;  
solve maximize <arithmetic expression>;  
solve minimize <arithmetic expression>;  
  
output [ <string expression>, ..., <string expression> ];
```

# Arrays

one- and multi-dimensional arrays which are declared using the type:

```
array [ <index-set-1>, ..., <index-set-n> ] of <type-inst>
```

They can contain: integers, enums, Booleans, floats or strings both as parameters and as variables (except for string)

Index sets are:

- ▶ an integer range,
- ▶ a set variable initialised to an integer range
- ▶ an enumeration type.

```
[ <expr-1>, ..., <expr-n> ] % 1d  
[ | <expr-1-1>, ..., <expr-1-n> |  
  ... |  
  <expr-m-1>, ..., <expr-m-n> | ]
```

```
array2d(1..3, 1..2, [1, 2, 3, 4, 5, 6])  
% is equivalent to  
[ | 1, 2 | 3, 4 | 5, 6 | ]
```

- ▶ concatenation operator ++, eg:  
[4000, 6] ++ [2000, 500, 500] evaluates to [4000, 6, 2000, 500, 500]
- ▶ a list is a one-dimensional array
- ▶ the built-in function length returns the length of a one-dimensional array.

# Array

A new example:

- ▶ modelling temperatures on a rectangular sheet of metal.
- ▶ approximate the temperatures across the sheet by breaking the sheet into a finite number of elements in a two-dimensional matrix.

```
set of int: HEIGHT = 0..h;  
set of int: CHEIGHT = 1..h-1;  
set of int: WIDTH = 0..w;  
set of int: CWIDTH = 1..w-1;  
array[HEIGHT,WIDTH] of var float: t; % temperature at point (i,j)  
  
% access as t[i,j]
```

- ▶ there is no explicit list type, one-dimensional arrays with an index set  $1..n$  behave like lists, we can use them as lists.

Laplace's equation states that when the plate reaches a steady state the temperature at each internal point is the average of its orthogonal neighbours.

```
% Laplace equation: each internal temp. is average of its neighbours  
constraint forall(i in CHEIGHT, j in CWIDTH)(  
    4.0*t[i,j] = t[i-1,j] + t[i,j-1] + t[i+1,j] + t[i,j+1])
```

On the edges the temperature is equal

```
constraint forall(i in CHEIGHT)(t[i,0] = left);  
constraint forall(i in CHEIGHT)(t[i,w] = right);  
constraint forall(j in CWIDTH)(t[0,j] = top);  
constraint forall(j in CWIDTH)(t[h,j] = bottom);
```

```
% corner constraints  
constraint t[0,0]=0.0;  
constraint t[0,w]=0.0;  
constraint t[h,0]=0.0;  
constraint t[h,w]=0.0;
```

Determine the temperatures in a plate broken into  $5 \times 5$  elements with left, right and bottom temperature 0 and top temperature 100

# Sets

For parameters: sets of integers, enums, floats or Booleans .

For variables: sets of integers or enums.

```
set of <type-inst> : <var-name> ;  
{ <expr-1>, ..., <expr-n> } % set  
<expr-1> .. <expr-2> % range
```

## Set operations

element membership	in
(non-strict) subset relationship	subset
(non-strict) superset relationship	superset
union	union
inter-section	intersect
set difference	diff
symmetric set difference	symdiff
number of elements in the set	card

# Enumerations

## Enumerated types enums

```
enum <var-name> ; % declaration as an unknown set of elements  
enum <var-name> = { <var-name-1>, ..., <var-name-n> } ; % definition via assignment
```

Elements of an enumerated type of  $n$  elements internally are represented as integers  $1 \dots n$ :

- ▶ they can be compared
- ▶ they are ordered, by the order they appear in the enumerated type definition,
- ▶ they can be iterated over,
- ▶ they can appear as indices of arrays,

They can be used for indexing

```
array[Products] of int: profit;
```

# List Comprehensions

Array and list comprehensions:

```
[i+j | i,j in 1..3 where j<i] % [2+1,3+1,3+2] = [3, 4, 5].  
{i + j | i, j in 1..3 where j < i} % {3, 4, 5}.
```

```
[ <expr> | <generator-exp> ]  
<generator>  
<generator> where <bool-exp> % filtering  
<identifier>, ..., <identifier> in <array-exp> % generator
```

Array comprehension can generate also variables while set comprehensions not.

# Built-in functions

**forall** takes a one-dimensional array and aggregate the elements.

The array is made of Boolean expressions (that is, constraints) and **forall** returns a single Boolean expression which is the logical conjunction

```
array of 1..3: a;  
forall( [a[i] != a[j] | i,j in 1..3 where i < j]);  
forall(i,j in 1..3 where i < j) ([a[i] != a[j]); % equivalent to the above one
```

The expressions mean:

```
a[1] != a[2] /\ a[1] != a[3] /\ a[2] != a[3]
```



# Aggregation Functions

Aggregation functions:

- ▶ for arithmetic arrays are: **sum**, **product**, **min**, **max**

When applied to an empty array, **sum** returns 0, **product** returns 1, **min** and **max** give a run-time error.

- ▶ for arrays containing Boolean expressions are: **forall** (logical conjunction, all hold), **exists** (logical disjunction, at least one holds), **xorall** (ensures that an odd number hold), **iffall** (an even number of holds).

Generator Call Expression:

```
<agg-func> ( [ <exp> | <generator-exp> ] )  
<agg-func> ( <generator-exp> ) ( <exp> )
```

(same syntax as **forall**)

# Production Planning Revisited

```
enum Products; % Products to be produced
array[Products] of int: profit; % profit per unit for each product
enum Resources; % Resources to be used
array[Resources] of int: capacity; % amount of each resource available
array[Products, Resources] of int: consumption; % to produce 1 unit of product

constraint assert(forall (r in Resources, p in Products)
  (consumption[p,r] >= 0), "Error: negative consumption");

int: mproducts = max (p in Products) % bound on number of Products
  (min (r in Resources where consumption[p,r] > 0)
    (capacity[r] div consumption[p,r]))

array[Products] of var 0..mproducts: produce; % Variables: how much should we make of each product
array[Resources] of var 0..max(capacity): used;

% Production cannot use more than the available Resources:
constraint forall (r in Resources) (
  used[r] = sum (p in Products)(consumption[p, r] * produce[p])
);
constraint forall (r in Resources) (
  used[r] <= capacity[r]
);
% Maximize profit
solve maximize sum (p in Products) (profit[p]*produce[p]);
output [ "\p) = \p(produce[p]);\n" | p in Products ] ++
  [ "\r) = \r(used[r]);\n" | r in Resources ];
```

## Data file for simple production planning model

```
Products = { BananaCake, ChocolateCake };  
profit = [400, 450]; % in cents  
Resources = { Flour, Banana, Sugar, Butter, Cocoa };  
capacity = [4000, 6, 2000, 500, 500];  
consumption= [| 250, 2, 75, 100, 0,  
              | 200, 0, 150, 150, 75 |];
```